PART/GNARL Interface Definition

Version 1.32

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1994/07/11 15:16:30

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This document describes the interface between the code generated by the Gnu NYU Ada Translator (GNAT) and the tasking support portion of the Gnu Ada Runtime Library (GNARL). The present version is a preliminary draft, for internal use of the GNAT and GNARL working groups, and for review by interested parties. We expect to make significant changes in both the content and scope of the document. Therefore, we warn against basing any product development on this version of this document.

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Section 1: Introduction

This document describes the interface between the code generated by the Gnu NYU Ada Translator (GNAT) and the tasking support portion of the Gnu Ada Runtime Library (GNARL). We call it the GNARL Interface (GNARLI, for short). This section describes why this document was produced, what it contains, and how it is intended to be used.

The interface between the GNARL and an underlying operating system or executive is defined in a separate document.

1.1 Purpose

The purpose of this document is to serve as a contract between the developers of an Ada compiler (GNAT) and the portion of an Ada runtime system (RTS) that supports tasking (the TRTS, for short).

1.2 Scope

The scope of the GNARLI is limited to the runtime system services that are required to implement the proposed Ada 9X [1] tasking semantics, and certain other runtime system services which may not be directly related to tasking but whose implementation we have not been able to separate cleanly from that of tasking.

The scope includes only RTS services that are invoked implicitly by compiler-generated code. In particular, it does not include those interfaces which are already defined by the Ada language, such as `Calendar` and `Text_IO`.

At present, the GNARLI covers the following services:

1. Task lifetime control
   (a) Task creation
   (b) Task activation
   (c) Task termination
   (d) `Terminated` attribute
2. Rendezvous management
   (a) Entry calls
   (b) **`Select`** and **`accept`** statements

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(c) `Callable` attribute
(d) `Count` attribute

3. Protected Objects
   (a) Protected Procedures and Functions
   (b) Protected Entry Calls
   (c) Count attribute
4. Abortion
5. Delays
6. Storage Management

Besides the Ada 9X features, we need to include support for interrupt entries.

1.3 Documentation Conventions

Subsections marked "RTS Responsibilities" describe features which an implementation of the GNARLI is required to provide. These are the features which the GNARLI client can count on.

Subsections marked "Compiler Responsibilities" describe invariants that the compiler or the code that it generates must maintain. Failure to do this results in unpredictable or erroneous behavior.

Subsections marked "RTS Notes" provides rationale for TRTS features, advice on the implementation the GNARLI, and other information on the intent of the interface. Text in these sections is not normative.

Subsections marked "Compiler Notes" provide advice on how GNARLI can be used to implement Ada features. Text in these sections is not normative.

All non-normative text is marked by grey areas in the margins to the left and right of the text. For example:

This paragraph is not normative.

A number of typographic conventions are also used to indicate the intended use of the associated text, as shown in the following list.

— `Ada program text`
— **`ada keywords`**
— *Definitions*

Change bars marked with the letter "a" indicate changes between this draft of the document and draft 1.28. a

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1.4 Principles

The intent of the GNARLI is to specify what the developer of a compiler needs to know about the TRTS and what the developer or tailorer of the TRTS needs to know about the compiler and the rest of the RTS. This involves specifying a clear division of duties.

There is clearly no single "best" way to divide these duties. The starting point for the GNARLI was CARTS [2], which was derived from MRTSI [5]. An attempt has been made to follow the CARTS interface faithfully, except where there is a very strong justification for difference. The main justifications for differences are:

1. Ada 9X requires new or different functionality.
2. A CARTS feature represents an extension to the Ada 83 tasking model, and we choose not to support that extension.
3. Design goals of the GNAT compiler conflict with some aspect of the CARTS interface.
4. CARTS operations provided for to support Ada 83 are redundant to operations needed to support Ada 9X.

In the design of extensions to support Ada 9X, an attempt has been made to follow the original principles of the CARTS and MRTSI interfaces, which were:

1. Necessity. Features should all be necessary to support standard Ada 9X for real-time applications.
2. Separation. The interface should cleanly separate the functions of the TRTS from those of the compiler-generated code and the rest of the RTS.
3. Information hiding. The interface should minimize the knowledge that the TRTS needs to have about the generated code and the rest of the RTS, and *vice versa*.
4. Implementation freedom. If possible, the interface should not force one particular implementation to the exclusion of another equally viable one.
5. Target-independence. Features should be specified in a way that is not specific to one machine or underlying operating system; that is, the interface should admit efficient implementations on a wide spectrum of machines.
6. Robustness. Features should favor the potential ability of the TRTS to detect and react gracefully to errors (if the implementor chooses), if this can be done without undue overhead.

These principles have several consequences. For example, a consequence of (1) and (2) is that none of the GNARLI features is intended to be invoked directly by an application programmer (though an implementation is not forbidden to export them

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to that level). Another consequence of (1) is that the GNARLI provides no support for debugging. An important consequence of (3) is that the GNARLI is procedural. The executive and compiler-generated code should not share access to any data structures—they communicate only through operations with explicit parameters.

1.5 Dependences of the TRTS on the Compiler

This document primarily describes the dependences of the compiler on the TRTS. Unfortunately, there are also some ways in which the TRTS implementation will (unavoidably) depend on the compiler. These include:

1. The representation of several types of data, including the predefined types `Integer`, `Boolean`, `Duration`, and `System.ADDRESS`, as well as all the types defined in the GNARLI.
2. Conventions for calling subprograms, passing parameters, and returning function values.
3. The representation of exceptions as data, and the mechanism of raising an exception (see Section 11).
4. The representation of a special exception for abortion. This exception cannot be caught by a **`when others`** exception handler. Since this exception is defined by the TRTS, it would not usually be visible to the application, and could not be caught by it.
5. The representation of task bodies, which must be compiled as procedures (see Section 6.7).

We assume the TRTS is written in Ada. Thus, most compiler-dependences will not be visible in the TRTS source code. This is the case with item 1, and, for the most part, item 2. However, there will remain some small parts of the TRTS implementation that will require specific knowledge of the compiler's conventions for subprogram calls. The required capability is described in this document as the TRTS calling a procedure "within the context of" a task (see Section 6.7). Examples of situations where the TRTS may need to do this are raising an exception in a blocked task, notifying a task of a hardware-detected exception, and (possibly) causing a task to abort itself. Sufficient information for the TRTS to be able to do this must be provided to the TRTS implementor by the compiler developer, and some (small, encapsulated) portion of the TRTS implementation will need to depend on it.

Item 3 shows up in TRTS interface via the exception ID parameter of `Exceptional_Complete_Rendezvous` and `Exceptional_Complete_Entry_Body` whose type must be visible to the TRTS.

Item 4 shows up in the TRTS interface in the `Abort_Signal` exception in the `System.Abortion` package (see Section 5). This document assumes that the compiler provides a `No_Others` pragma that takes an exception as an argument, and which prevents that exception from being caught by a **`when others`** handler.

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The remaining dependences of the TRTS on the compiler are not visible in the TRTS interface, but some will be manifest in the form of small segments of compiler-dependent code within the TRTS.

1.6 Dependences of the Compiler on the TRTS

The `System.Machine_Specifics` package contains declarations of types that represent conventions about the use of the target machine architecture on which the compiler, TRTS, and other runtime system components may depend. The structure, size, representation, and contents of these types would ordinarily be determined by the compiler developer, based on what is natural and efficient for the specific target machine. The TRTS should be coded in a modular fashion, so that portions which depend on the definitions of these types can be modified easily. Moreover, for a given target machine, it should be possible to fix a set of definitions for these types which can be common across compilers for that machine.

1.7 Dependences of the TRTS on the Target Machine

```
package System.Machine_Specifics is
   type Init_State is implementation-defined;
   type Pre_Call_State is private;
   type Task_Storage_Size is range implementation-defined;
   type Machine_Exceptions is range implementation-defined;
   type Interrupt_ID is range implementation-defined;
   type Interrupt_Info is implementation-defined;
   type Error_Information is implementation-defined;
private
  ... implementation-defined
end System.Machine_Specifics;
```

1.7.1 Initial Task State

The `Init_State` type shall be used to represent the initial hardware state of a task, when it starts executing. It contains information needed for initialization of a task. It must allow the code for the body of the task to be executed within a new thread of control, and allow the Ada environment for this code (e.g. uplevel references) to be established.

1.7.1.1 GNARL Implementation Notes

The specific details below are preliminary specifications for the GNAT Ada compiler and GNARL TRTS for the Sun SPARC[1] architecture.

Ada 9X provides subprogram access values that allow the referenced subprogram to be called in the proper environment. One way to implement `Init_State` values is to compile task bodies as procedures and use a procedure access type to these

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subprograms as the implementation of `Init_State`. A task could then call this procedure through an `Init_State` value at activation.

Whether this works or not depends on the implementation of uplevel references. The Ada 9X reference manual 2.0 3.10(7.a) points out that for implementations which use a global display, it will generally be possible to call a subprogram using only the address of the subprogram code, since the environment needed for uplevel references will be in the display at the point of call. On such implementations, it may also be possible for the TRTS to call the task body procedure from the new thread of control using only the address of its code, but this will depend on how "global" accesses are done. If some context is needed to access the global display (e.g. there is a distinguished register pointing to global data), then `Init_State` may have to be defined to allow this context to be established in the new thread.

In systems that use static chains for uplevel references, it will usually be desirable to include the address of the statically enclosing frame with the code address as part of a subprogram access value[2]. In this case, the TRTS can use this value to call the task procedure in the new thread of control using the normal Ada calling mechanism.

The GNAT compiler does implement access to subprogram in this way, so `Init_State` will be implemented as an access to procedure type in GNARL:

```
private
   type Init_State is access procedure(Arg: System.Address);
   ... implementation-defined
end System.Machine_Specifics;
```

Note that `Init_State` is a global procedure access type. Ada accessibility rules only allow access values referring to library level procedures to be assigned to variables of such a type. This is to prevent dangling references in the form of access variables referring to procedures whose static environment no longer exists. Using access values to execute task body procedures depends on those procedures being nested at the level of the task bodies that they represent, so in general it will violate Ada semantics to assign access values referring to these procedures to variables of type `Init_State`. Note that this conservative compile-time check is not needed in this case, since Ada task termination processing assures that the environment of a task will exist until the task is terminated. There must therefore be some means of circumventing accessibility checks to allow arbitrary procedure pointers to be assigned to variables of `Init_State`; for GNAT, this consists of using `Unchecked_Conversion` to convert procedure access values to `Init_State` values.

1.7.2 Pre-Call State

The `Pre_Call_State` type shall be the portion of the hardware state of a task that is not preserved by the subprogram-calling convention. For example, this might be

[2] Even in this case, it is possible to do without a static link in a procedure access value. A subprogram can find the enclosing frame from the current frame by traversing the static links to the outermost frame and then traversing them back down for the requisite number of nesting levels. However, it is unlikely that any implementation would choose to add this much distributed overhead to subprogram calls to reduce the size of access to subprogram values by one address.

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an array of integers, representing the values of certain registers. (No relationship is specified between this type and the `Init_State` type.)

1.7.3 Task Storage Size Type

The `Task_Storage_Size` type shall be used to specify task storage-sizes. This shall be an integer type.

1.7.4 Machine Exceptions Type

The `Machine_Exceptions` type shall be used to identify a hardware-detected exception when notifying the exception-recovery system in the rest of the RTS. This might be an enumerated or integer type.

1.7.5 Interrupt ID Type

The `Interrupt_ID` type shall be used to identify hardware interrupts and virtual interrupts.

Need to decide whether to merge POSIX signals with interrupts, or provide a special type and package(s) for signals.

1.7.6 Interrupt Info Type

The `Interrupt_Info` type shall be used to specify information (if any) that is passed from a hardware interrupt to an interrupt handler procedure.

1.7.7 Error Information Type

The `Error_Information` type shall be used to provide any additional information associated with a hardware-detected error. For example, this might be a record type.

1.8 RTS Service Call Conventions

TRTS actions are assumed to be initiated by one of the following:

1. Request by an Ada task (including the environment task), which is formally termed a "TRTS call". From the point of view of the calling task, these are indistinguishable from ordinary subprogram calls.
2. A hardware interrupt (or trap). Examples of this include completion of an I/O operation, and expiration of a delay.
3. A software signal. These are similar to interrupts or traps, but are generated by an underlying operating system. They may serve as an interface to interrupts and traps, or exist together with them.

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Only TRTS calls are visible at the TRTS interface. Interrupt, trap, and signal handlers are installed by the TRTS, and are therefore entirely internal to the TRTS implementation.

The syntax of Ada subprograms is used to name the TRTS calls and specify their parameters.

The TRTS is assumed to be written in Ada. Therefore, the compiler should be able to generate in-line code for those TRTS calls where it makes sense. Ideally, the compiler may perform optimizations on the expanded TRTS code.

The compiler is expected to provide correct arguments to the TRTS. In particular, if there is a constraint on a parameter to a TRTS call the compiler is expected to perform any necessary checking and raising of exceptions, or ensure that the constraint is always satisfied.

The TRTS may (but need not) check for bad arguments. If such an error is detected, the TRTS is free to choose an appropriate response.

1.9 Uninterpreted Data in the TRTS

Some TRTS operations require the TRTS to transfer *uninterpreted data* between the compiler-generated code of one task and that of another.

1.9.1 RTS Responsibilities

When uninterpreted data must be passed between tasks, the TRTS shall provide a means of transferring a a single value of type `System.Address` from one task to the other. The TRTS shall provide the means for task identification and synchronization necessary for this transfer, as documented under the individual operations.

1.9.2 RTS Notes

An example of where uninterpreted data is needed is task activation. The value of these discriminants are known to the task's activator, but must be available within the task itself. The TRTS must perform this transfer, since it mediates communication between tasks. However, the TRTS itself does not need to access the discriminants, so the TRTS simply accepts uninterpreted data when creating a task and later provides this data to the new task.

1.9.3 Compiler Responsibilities

Exactly how the compiler-generated code uses a `System.Address` value specifying uninterpreted data is up to the compiler.

1.9.4 Compiler Notes

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Compiler generated code should assure that the `System.Address` parameter can be used to access the appropriate uninterpreted data. All tasks are assumed to share a common virtual address space, so this address can be used as a pointer to a data block, allowing access to arbitrary data.

1.10 Executing Compiler-Generated Code

The TRTS effects the execution of compiler generated code in complex ways. This is more true of Ada 9X than it was with Ada 83, with the addition of complex control structures such as asynchronous **`select`** statements and protected objects. One way to allow the TRTS to control execution is to pass pointers to compiler-generated subprograms into the TRTS. These subprograms execute fragments of the structures whose execution must be controlled by the TRTS. The TRTS can then call these subprograms (a *callback* to compiler-generated code) as appropriate.

Another way for the TRTS to control execution is through embedded calls to the TRTS within compiler-generated code. This gives the TRTS the opportunity to block the task or communicate with the compiler-generated code through parameters. This only allows for local control of execution; the influence of the TRTS is to code immediately preceding or following the call, in general requiring support from additional compiler-generated code. This makes for a more complex interface and for more complex compiler-generated code to use it. However, it has the advantage of greater isolation of the TRTS and the compiler-generated code, allowing greater protection between the two, and requiring less protection within the TRTS for callbacks to less predictable application code. For these reasons, the GNARLI avoids the use of callbacks.

1.11 Package Structure

The TRTS interface is described via several Ada package specification, based on a natural functional division. These packages are:

1. `System.Task_IDs`—in Section 4.
2. `System.Abortion`—in Section 5.
3. `System.Task_Stages`—in Section 6.
4. `System.Delays`—in Section 7.
5. `System.Rendezvous`—in Section 8.
6. `System.Protected_Objects`—in Section 9.
7. `System.Storage_Allocation`—in Section 10.

Both the specification and body of any of the GNARLI packages may include additional declarations of types, objects, and operations required by the implementation

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of that package or other TRTS packages. The compiler must not rely on the existence (or nonexistence) of any such declarations.

The GNARLI package specifications are presented more fully in the following sections, with the intended semantics and rationale for each of their individual components. In some cases the description of the semantics refers to the Ada language definition.

1.12 Examples of GNARLI Use

To clarify the intended use of the GNARLI, this document provides examples of Ada tasks implemented using it. This is usually done by showing an Ada code fragment containing task structures, followed by a translation of this fragment into Ada code free of tasking features using GNARLI constructs. Please note that these examples are to be understood in a conceptual sense, for illustrative purposes only. In the actual implementation, all Ada constructs almost certainly would be translated directly to machine code by the compiler, rather than back into Ada code, and compiler optimizations are likely to result in some simplifications.

In fact, it is not always possible to provide the correct tasking semantics with Ada code. For this reason, some examples are expressed using extensions to the Ada language. These are defined as needed in the text accompanying the examples.

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Section 2: Error Checking at the TRTS Interface

This section describes general rules regarding errors, and propagation of exceptions, at the TRTS interface.

2.1 RTS Responsibilities

TRTS calls shall not propagate exceptions back to the caller, except as explicitly required or permitted in this document.

The TRTS is responsible for causing certain correct TRTS calls to propagate a language-defined exception, under certain conditions. For example, an attempt to rendezvous with a terminated task causes `Tasking_Error` to be raised. This is considered part of the "normal" behavior of the TRTS call, and is specified as such in the description of the RTS responsibilities for the TRTS call.

In contrast, there are possible situations where a call to the TRTS violates the conventions of the TRTS interface. The compiler is trusted to produce correct TRTS calls, so such a violation is a compiler defect. The TRTS is *not required* to check for incorrect TRTS calls. The effect of an incorrect TRTS call is unspecified by the GNARLI.

Elaboration of the TRTS packages (specifications or bodies) shall raise no exceptions. In particular, elaboration of the TRTS packages shall not raise `Program_Error`, due to access before elaboration problems in the TRTS implementation.

2.2 RTS Notes

If a TRTS implementation does not check for incorrect TRTS calls, the result of an incorrect call may be to "crash" the system. The TRTS implementor may choose to check for incorrect calls, to preserve the integrity of the TRTS. If so, the choice of error recovery action may be made, *ad hoc*, by the TRTS implementor.

2.3 Compiler Responsibilities

The compiler is responsible for generating only correct TRTS calls. In particular, all parameter values must be valid, and satisfy any constraints imposed by this document.

2.4 Compiler Notes

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The compiler should generate TRTS calls in a fashion that if the TRTS performs its specified responsibilities the result will be consistent with the Ada language definition. However, the GNARLI implementation should not attempt to check for errors in implementing Ada semantics.

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Section 3: Predefined Packages

This section describes aspects of the TRTS interface related to the packages `System` and `Standard`.

3.1 System Package

3.1.1 RTS Responsibilities

The TRTS implementation shall use the mechanism provided by the compiler implementor to ensure the constants and subtypes in the package `System` are consistent with the ranges of priorities and the clock resolution supported by the TRTS.

3.1.2 RTS Notes

The range of priorities used internally by the TRTS may be wider than that provided to the Ada user. However, it cannot be narrower, since tasks with distinct Ada priorities must have distinct priorities in the TRTS, if the Ada 9X Real-Time Annex is supported.

3.1.3 Compiler Responsibilities

The compiler implementor is responsible for providing a mechanism that permits the TRTS implementor to define the values and ranges of the following constants and subtypes, which are declared in the package `System`:

1. `Tick`, the nominal tick of `Calendar.Clock`.
2. `Priority`, the range of non-interrupt priorities which is supported.
3. `Any_Priority`, the full range of priorities supported.
4. `Interrupt_Priority`, the range of interrupt priorities supported.

The TRTS implementation must be able to define child packages of `System`. In particular, it must be able to define those required by the Ada 9X Annexes, including `System.Real_Time`. All TRTS packages are declared as children of package `System`.

The TRTS implementation packages (specifications as well as bodies) must be allowed to depend on the package `System`, as well as child packages defined by the RTS.

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3.2 Standard Package

3.2.1 RTS Responsibilities

The TRTS implementation shall use the mechanism provided by the compiler implementor to ensure the type `Duration` is consistent with the clock resolution supported by the TRTS.

3.2.2 RTS Notes

The representation of the standard `Duration` type is presumed to be built into the compiler. However, for efficient conversion, it is required that the time unit used in the implementation of delays be an integral multiple of `Duration'Small`.

3.2.3 Compiler Responsibilities

The compiler implementor is responsible for providing a mechanism that permits the TRTS implementor to tailor the value of `Duration'Small`, so that it is an integral multiple or divisor of the underlying clock tick.

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Section 4: Representing Tasks

This package introduces the `Task_ID` type, which is used to represent tasks at the TRTS interface. It also provides some fundamental operations on task IDs.

4.1 Package Specification

```
package System.Task_IDs is
   type Task_ID is private;
   Null_Task: constant Task_ID;
   function Same_Task(L, R: Task_ID) return Boolean;
   function Self return Task_ID;
   subtype  Task_String is String(1..16);
   function Image(T: Task_ID) return Task_String;
private
   type Task_ID is ...;
end System.Task_IDs;
```

4.2 Representing Tasks to the TRTS

4.3 RTS Responsibilities

A value of the `Task_ID` type (a *task ID*) is *valid* if and only if it identifies some task that has been created by the `System.Task_Stages.Create_Task` and the corresponding task has not subsequently been expunged.

A `Task_ID` value that is returned by the `System.Task_Stages.Create_Task` operation represents the created task to the TRTS (here distinguished from an unelaborated task declaration, task type, or unevaluated task allocator) until the task is expunged.

Note that the value `Null_Task` is not a "valid" task ID, though this value shall be accepted by certain GNARLI operations with parameters of the `Task_ID` type.

A task shall be *expunged* by the TRTS (only) at the earliest of the following:

1. Exit from the task's immediate master (including termination of the master, if the master is a task).

2. A call to `Expunge_Unactivated_Tasks` for the activation list containing the task.

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The `Same_Task` function shall return True if-and-only if the `L` and `R` parameters specify the same task or are both equal to `Null_Task`.

`Self` shall return a value of type `Task_ID` that uniquely identifies the calling task.

`Image` shall return a string that uniquely identifies the task identified by the `T` parameter, or that identifies the null task if `T` is equal to `Null_Task`. The effect is undefined if `T` is invalid and is not equal to `Null_Task`.

4.4 RTS Notes

The TRTS must be able to determine the identity of the currently executing Ada task on a processor (in the case of an interrupt handler, it also needs to be able to determine the identity of the interrupted task.) This operation must be implemented very efficiently, perhaps as in-line code, since it is likely to be used frequently. It generally requires some cooperation with the compiler. One solution is to reserve a global register to contain a reference to the currently executing task. A variation on this is to store the task ID at one end of the corresponding task's runtime stack, or to put a pointer to it at that location.

It is recommended that the TRTS implementation provide a mechanism for using the string returned by `Image` to determine the declaration of the corresponding task or task type. Since the `Image` function is only provided for debugging output, the mechanism need not be available at run time in the target environment. (See also the `Image` function for the `Exception_ID` type, Section 11.1.)

4.5 Compiler Responsibilities

For TRTS calls which require a valid task ID as parameter, the compiler is responsible for ensuring that it is valid.

4.6 Compiler Notes

A valid task ID represents a unique Ada task. However, there may be more than one task ID value which represents a given task. For this reason, the `Same_Task` function is provided. This operation should always be used to determine if two `Task_ID` values represent the same task. a

The compiler-generated code may need to access the task ID corresponding to the currently executing task. For this reason, the `Self` operation, already needed by the TRTS, is made visible in the interface.

When the name of a task type is used to refer to the object that designates the task currently executing the body of the task type (see [1], Section 9.1(15)), the compiler must find the corresponding `Task_ID` value. The function `Self` can be used directly, but it will generally be more efficient to call this function once whenever a task of the type in question starts to activate. This would probably be done by keeping the task

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ID in a local variable of the task and translating the name of the task type within the body of the task to a reference to this local variable. This will work even if the reference to the task comes from some other task nested within it.

An task ID is no longer valid after return from a subprogram where the task is declared. It is the compiler's responsibility to raise `Program_Error` in functions that attempt to return local tasks, since tasks are return-by-reference types.

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Section 5: Abortion

The System.Abortion package provides operations to abort tasks or sequences of statements (for ATC).

5.1 Package Specification

```
with System.Task_IDs;
package System.Abortion is
   Max_ATC_Nesting: constant Natural:= implementation-defined;
   Abort_Signal: exception;  -- pseudo-exception
     pragma No_Others(Abort_Signal); -- cannot be caught by when others =>
   procedure Defer_Abortion;
   procedure Undefer_Abortion;
   type Task_List is array(Positive range <>) of Task_IDs.Task_ID;
   procedure Abort_Tasks(Tasks: Task_List);
end System.Abortion;
```

The description of this package is divided into several subsections.

5.2 ATC Levels

Asynchronous **select** statements can be nested, so that a task can have several entry calls pending at once. Max_ATC_Nesting is the maximum number of ATCs that can be pending on a given task at one time.

5.2.1 RTS Responsibilities

The TRTS shall allow a task to make a combined total of Max_ATC_Nesting calls to the Task_Entry_Call or Protected_Entry_Call operations that can be pending on a task simultaneously. The effect of attempting to queue more than Max_ATC_Nesting entry calls is undefined.

If Max_ATC_Nesting is zero, the implementation places no fixed limit on the number of entry calls that can be pending on a task at one time. The number of entry calls is limited only by the availability of the necessary TRTS resources in this case.

5.2.2 RTS Notes

In Ada 83, a task was always blocked when it made an entry call, so only one such call could be pending on a task at a time. This could be implemented by a single queue element, possibly part of a task control block. In contrast, an asynchronous **select** statement may be queued waiting for an entry while the calling task continues to

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execute, and this can be repeated any number of times. This implies some scheme for dynamically allocating queue elements to allow a task to be on multiple queues, subject only to storage availability.

Allowing the implementation to limit the depth of nesting allows for the preallocation of all needed queue blocks at task creation time, avoiding the need to do dynamic allocation as part of the entry call mechanism. This preallocation could take the form of `Max_ATC_Nesting` queue elements allocation as part of a task control block.

5.3 The Abort Signal Exception

Task abortion requires the abortion of the sequence of statements of the task body; ATC requires the abortion of the sequence of statements comprising the abortable part. This interface uses a special exception to perform abortion.

5.3.1 RTS Responsibilities

The RTS will defer abortion and raise the `Abort_Signal` exception asynchronously in a task whenever it is aborted or completes the triggering statement of an asynchronous select statement.

The `Abort_Signal` exception is raised by various TRTS calls responsible for completing triggering statements, including `Complete_Rendezvous` (Section 8) and `Complete_Entry_Body` (Section 9).

Abortion will remain deferred until the abortable part or task being aborted is left [1].

5.3.2 RTS Notes

Task abortion and ATC both cause a task to stop its current execution no later than an abortion completion point. A task that has been aborted will not execute again once the abortion has taken effect, where ATC will resume execution after completing some sequences of statements. In both cases, all finalization for the completed structures must be performed.

This similarity makes it desirable to use the same mechanism for both features. This interface allows a task to be aborted as if it were the abortable part to some enclosing asynchronous select statement.

5.3.3 Compiler Responsibilities

The compiler must support a special kind of pseudo-exception, `Abort_Signal`, which cannot be caught by a handler for **`others`**. A possible interface for this is indicated above by the pragma `No_Others`.

[1] The GNAT compiler is designed to defer abortion between the point at which an exception is raised until it is propagated. This effects the deferral of abortion during the propagation of `Abort_Signal` without further intervention by the GNARL. a

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5.4 Deferring Abortion

The `Defer_Abortion` and `Undefer_Abortion` operations allow the effect of abortion on a task to be controlled.

5.4.1 RTS Responsibilities

`Defer_Abortion` shall defer the effects of abortion in the calling task. The TRTS shall keep track of all abortion deferrals, either by `Defer_Abortion` or by some other means (e.g. a call to `Lock`). It shall continue to defer abortion until all such deferrals are removed. Abortion itself also defers abortion, and as such must be completed by some TRTS call that undefers abortion.

`Undefer_Abortion` shall raise the `Abort_Signal` exception if it corresponds to the earliest pending `Defer_Abortion` call made by this task, and it finds that abortion to an outer asynchronous call level or the whole task is pending. Otherwise it will remove the effects of the most recent call to `Defer_Abortion` made by the current task. Note that `Undefer_Abortion` leaves abortion deferred if it raises the `Abort_Signal` exception; abortion is always deferred while abortion is in progress.

The effect of calling `Undefer_Abortion` is undefined unless abortion is deferred. The effect of calling `Undefer_Abortion` when abortion has been deferred most recently in the calling task by a call other than `Defer_Abortion` (e.g. `Lock`) is undefined.

5.4.2 RTS Notes

The intent of the definition of `Defer_Abortion` and `Undefer_Abortion` is that abortion deferral be nestable. For example, a `Defer_Abortion`/`Undefer_Abortion` pair can be called within a protected object without concern that the `Undefer_Abortion` will allow abortion within an eggshell, which would violate Ada semantics.

5.4.3 Compiler Responsibilities

The compiler shall not call `Undefer_Abortion` unless abortion is deferred, and only if the most recent abortion deferral was the result of a call to `Defer_Abortion`.

5.4.4 Compiler Notes

Abortion should be deferred during the execution of certain constructs, including protected actions, finalizations, while waiting for the completion of a rendezvous in progress, and during abortion itself. In most cases, abortion can be controlled as part of the TRTS calls used to implement these constructs. However, there are vulnerable sections of code that are under the control of the compiler, such as finalization, where the compiler must assure that abortion is deferred (see Section 6.5.3).

The `Abort_Signal` exception should be handled in each nested asynchronous select statement being executed by the task in LIFO order, making the appropriate calls to

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the TRTS to undefer abortion. If the asynchronous **select** statement is nested in another construct that is being aborted, this will raise the `Abort_Signal` exception again; otherwise it will undefer abortion and return. This in effect allows the TRTS to decide when the correct level of ATC has been aborted.

5.5 Aborting Tasks

The `Abort_Tasks` operation allows a group of tasks to be aborted.

5.5.1 RTS Responsibilities

`Abort_Tasks` shall abort the set of tasks specified by the `Tasks` parameter. If a task is suspended and not abortion deferred, it will be completed before `Abort_Tasks` returns. Tasks dependent on the aborted tasks are also aborted. It also causes all operations that undefer abortion (`Undefer_Abortion`, above; `Unlock`, in Section 9) to re-raise the exception whenever they are called by the affected task.

The effect of `Abort_Tasks` is undefined if `Tasks` contains any references to invalid task IDs or the lower index bound is not 1.

5.5.2 Compiler Responsibilities

The parameter `Tasks` of `Abort_Tasks` shall be a list of valid task IDs. The lower index bound of the array shall be 1.

5.5.3 Compiler Notes

The **abort** statement corresponds directly to the `Abort_Tasks` procedure. This will raise the abortion exception in the affected tasks. Whenever the `Abortion` exception is handled (as for an asynchronous **select**) statement, and attempt to undefer abortion should be made. This raises `Abort_Signal` again if the abortion is to a higher level, which will always be the case when the whole task is aborted. Eventually, target tasks will reach the abortion handler for the whole task, which will call `Complete_Task`. The task procedure will then return to the TRTS, which will terminate it (see Section 6.12.4).

The statement `Abort_Tasks((1=>T1,2=>T2));` is not necessarily equivalent to `Abort_Tasks((1=>T1)); Abort_Tasks((1=>T2));`. The former call will abort `T2` even if `T1` is the caller or is dependent on the caller, whereas `Abort_Tasks((1=>T1))` will never return in this case, and will therefore never call `Abort_Tasks((1=>T2))` to abort `T2`.

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Section 6: Task Creation and Activation

This section describes the TRTS interface that supports transitions between the stages of a task's lifetime, from nonexistence through activation, normal execution, completion, and termination. This interface is provided by the package `System.Task_Stages`.

6.1 Package Specification

```
with System.Machine_Specifics;
with System.Rendezvous;
with System.Task_IDs; use System.Task_IDs;
package System.Task_Stages is
  type Master_ID is private;
  function  Current_Master return Master_ID;
  procedure Enter_Master;
  procedure Complete_Master;
  type Access_Boolean is access Boolean;
  type Size_Type is new Machine_Specifics.Task_Storage_Size;
  Unspecified_Size: constant Size_Type:= Size_Type'First;
  type Activation_Chain is limited private;
type Activation_Chain_Access is access all Activation_Chain;                    |a
  Unspecified_Priority: constant Integer:= implementation-defined;
  procedure Create_Task(
      Size: Size_Type;
      Priority: Integer;
      Num_Entries: System.Rendezvous.Task_Entry_Index;
      Master: Master_ID;
      State: Machine_Specifics.Init_State;
      Uninterpreted_Data: System.Address;
      Elaborated: Access_Boolean;
      Chain: in out Activation_Chain;
      Created_Task: out Task_ID);
procedure Activate_Tasks(                                                       |
      Chain: Activation_Chain_Access);                                          |a
  procedure Expunge_Unactivated_Tasks(Chain: in out
Activation_Chain);
  procedure Complete_Activation;
  procedure Complete_Task;
  function  Terminated(T: Task_ID) return Boolean;
private
  type Master_ID is ...;
  type Activation_Chain is ...;
end System.Task_Stages;
```

The description of this package is divided into several subsections.

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6.2 Task Masters

6.2.1 RTS Responsibilities

A value of the type `Master_ID` is called a *master ID*. A master ID shall be *valid for a task* `T` if and only if it has been returned by a call to `Current_Master` from `T`, and only until `T` executes a call to `Complete_Master` or `Complete_Task`.

A valid master ID shall uniquely identify a master within a particular task and its dependents (here distinguished from other tasks and masters within them) until that master is exited.

6.2.2 Compiler Responsibilities

The compiler shall assure that a stack frame is generated for each nontrivial master.

6.2.3 Compiler Notes

The concept of a "master" is the basis of the rules for task termination in Ada. A master is a construct that performs finalization of local objects after it is complete, but before leaving, and can be a task body, block statement, subprogram body, or **`accept`** statement.

Task objects declared by the master, or referred to by objects of access types declared by the master, are said to *depend* on that master. Dependent tasks must terminate before a master performs finalization on other objects that it has created. Since **`accept`** statements have no declarative part, tasks are never dependent on them, and so a master for purposes of task termination is a task body, block statement, or subprogram body.

Finalization of objects other than tasks is the responsibility of the compiler, and is not part of the TRTS. The TRTS is only concerned with masters for purposes of task termination. For this reason, a master which is a task body or has a local declaration of either (1) a task object or an object with a task component, or (2) an access type whose designated type is a task type or a type with a task component is referred to as a *nontrivial master*; in the exposition that follows, the term master will generally mean a non-trivial master. In order to properly implement task termination, the TRTS must be notified of the change in master-nesting whenever execution is about to enter or leave a nontrivial master. This is done via procedures `Enter_Master` and `Complete_Master` (or, in the case of tasks, via `Create_Task` and `Complete_Task`), which are described further below.

Master IDs are used to specify masters of tasks. A master ID is guaranteed to be valid for a task so long as it is only referenced from within the master which obtained that value by calling `Current_Master` (including masters that may be called from or nested within that master).

The requirement for a stack frame for each nontrivial master is to allow the TRTS to use the combination of current frame pointer and task ID for the current task to

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determine the master of a task (see Section 6.4.2, below).

6.3 Current Master

The `Current_Master` function returns a value of the `Master_ID` type that identifies the master from which the function is called.

6.3.1 RTS Responsibilities

A master ID that is returned by the `Current_Master` function shall uniquely identify the currently executing master within the calling task until that master is exited. A master that is not a task shall be considered to have been exited when the `Complete_Master` operation is called by that master, and a master that is a task shall be considered to have been exited when the `Complete_Task` operation is called by that task.

6.3.2 Compiler Notes

`Current_Master` should be used by the compiler-generated code to associate a master ID with tasks dependent on it. For task object declarations, this can be done by calling `Current_Master` at the point where the task is created (see `Create_Task`, below). For task access types, `Current_Master` should be called at the point where the type is elaborated and saved, so that it can be accessed for any task created by the allocation of an object of that type.

6.4 Master Entry

`Enter_Master` notifies the TRTS that execution is entering a new master.

6.4.1 RTS Responsibilities

The TRTS is responsible for keeping track of the master nesting of each task.

6.4.2 RTS Notes

If the TRTS keeps track of master nesting via a counter per task, this procedure can be implemented rather simply, as an increment operation on the master nesting count of the calling task.

Another possibility is to use the master frame pointer as the level counter. With this implementation, `Enter_Master` does not perform any operation. `Current_Master` returns the current frame pointer, which the compiler-generated code associates, together with the *parent* task (the task which elaborated the master), with the tasks dependent on the master. In fact, the frame pointer alone is adequate to identify the master, since only the parent's stack can contain a given master frame pointer. However, explicitly associating the parent with each dependent task avoids the need to search all tasks for it.

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When `Complete_Master` is called, the frame pointer of the master at that point is compared with the master ID associated with children of the task that elaborated the master. Children which are more deeply nested must terminate before the master can complete. This search can be made less expensive by having each task keep a list of its children. When one of these masters is exited, only the list of the task which elaborated the master needs to be searched.

6.4.3 Compiler Responsibilities

The compiler shall ensure that `Enter_Master` is called (exactly once) whenever execution enters a nontrivial master. It shall not be called on entry to task masters that are task bodies. (For masters that are task bodies, the effect is subsumed by `Create_Task` and for the environment task the effect is subsumed by start-up of the TRTS of the associated active partition.)

6.4.4 Compiler Notes

The compiler is permitted to call `Enter_Master` on entry to a trivial master. This may be desirable in cases where the compiler cannot determine if a master is trivial or not. However, the implementation of `Enter_Master` may involve significant overhead, so the compiler should avoid calling it for trivial masters whenever possible.

6.5 Master Exit

`Complete_Master` notifies the TRTS of completion of a master.

6.5.1 RTS Responsibilities

For this operation, the TRTS shall perform all actions required before exit from a master of tasks. These actions are:

1. Terminate any dependent tasks that have not been activated.
2. Block the task calling `Complete_Master` until all the tasks dependent on the current master are either terminated or are waiting on an open **`terminate`** alternative. Once this state has been reached, terminate the tasks waiting on **`terminate`** alternatives.
3. When all the tasks dependent on the master have terminated, deallocate any remaining storage associated with these tasks.
4. Update the TRTS's internal state to recognize the (new) currently executing master's identity within the task that made the TRTS call.

The effect of `Complete_Master` is undefined if there is a rendezvous that was begun by a call to `Selective_Wait` or `Accept_Simple` from within that master but has not yet been completed by a matching call to `Complete_Rendezvous` or `Exceptional_Complete_Rendezvous`. It is also undefined if called from a task executing a protected action.

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6.5.2 Compiler Responsibilities

The compiler shall ensure that `Complete_Master` is called (once) after the master completes, whether normally, via a **`return`**, **`exit`** or **`goto`** statement, or through propagation of an exception, but prior to finalization of the master. It shall not be called on exit from masters that are task bodies. (The effect in the latter case is subsumed by `Complete_Task`.)

After return from a call to `Complete_Master`, or termination of a task that is a master, no further reference to this master's ID or to the task ID of any dependent task of this master shall be made by the compiler-generated code. (Such master IDs will no longer be valid. Of course, all such dependent tasks must have terminated by this time.)

6.5.3 Compiler Notes

`Complete_Master` must be called before finalization of a task master, regardless of how the master was completed. This requires that `Complete_Master` be called in a handler that is always executed, and which is protected against abortion. Even if the master is exited due to an exception raised in its declarations, it must call `Complete_Master` to terminate tasks that have been created but not yet activated.

This requirement on the compiler is beyond that made by the Ada language itself for exception handling and finalization. For example, consider a declare block which contains tasks:

```
declare
  ... mixed task and other object declarations
begin
  ... code for the sequence of statements of the master
end;
```

A possible translation of this to use TRTS calls might be:

```
declare
  Enter_Master;
  ... code for elaboration of master declarations
begin
  ... code for the sequence of statements of the master.
when others=>
  Complete_Master;
  raise;
end;
```

However, if an exception is raised in the elaboration code, it will be propagated; exception handlers are not executed unless the exception is raised in the corresponding handled sequence of statements, which does not have a declarative part. This would mean that `Complete_Master` would not be called. `Complete_Master` cannot be called from a dynamically enclosing exception handler, since the finalization of the declare block would already have been done, contrary to Ada semantics.

A possible alternative is to use the finalization mechanism to call `Complete_Master`.

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For illustrative purposes task completion is shown implemented by a dummy controlled object created by the last declaration in the procedure representing the task body. The idea is that, since finalization is performed in LIFO order, `Complete_Master` will be called before any user finalization is performed:

```
Complete_Master_Type is new Finalization.Root_Limited_Controlled;
procedure Finalize(C: in out Complete_Master_Type) is
begin
  Complete_Master;
end Finalize;
declare
  Enter_Master;
  ... mixed master and other object declarations
  Complete_Master_Object: Complete_Master_Type;
begin
  ... code for the sequence of statements of the master.
end;
```

However, this also fails if an exception is raised during the declarative part of the master, since `Complete_Master_Object` will not have been elaborated yet, and will therefore not be finalized.

Any compiler will have to solve this problem with an extension to Ada semantics. The solution used in the GNAT compiler is a special clause that can be added to any master containing a sequence of statements that will be executed when that master (in the Ada sense, including **accept** statements) is left, either through normal completion or through the propagation of an exception (including the `Abort_Signal` exception), even when the exception is raised in the declarative part. The syntax that has been adapted to express this construct is a sequence of statements, starting with the keywords **at end**, between the exception handlers (if any) and the **end** keyword of the master.

```
declare
  Enter_Master;
  ... mixed master and other object declarations
begin
  ... code for the sequence of statements of the master.
at end
  Complete_Master;
end;
```

The **at end** clause will be used throughout this document to express operations that must be executed when leaving a master or other construct which requires cleanup before being left. If the handler was entered as the result of an exception, the exception is propagated to the surrounding execution when the handler completes.

An additional problem is posed by asynchronous ATC. If the `Abort_Signal` exception can be raised at any point in the code, any handler is vulnerable to being aborted before it can complete its processing. For TRTS calls like `Complete_Master` which must complete to ensure proper Ada semantics, this cannot be permitted. Note that it is not enough to simply defer abortion in `Complete_Master`, since the handler might be aborted before it is called.

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For this reason, the compiler assures that abortion is deferred before the handler is entered, and that it is enabled again by calling `Undefer_Abortion` when the handler is complete. If the handler is entered because of an `Abort_Signal` exception, abortion will already be deferred. In other cases, the TRTS must defer it using `Defer_Abortion`, perhaps as part of raising the exception.

The `Undefer_Abortion` operation at the end of the **`at end`** handler will either undefer abortion or raise the `Abort_Signal` exception, leaving abortion deferred. The propagation of an `Abort_Signal` exception indicates that abortion is in progress, and it should not be propagated unless abortion is deferred. Therefore, unlike other exceptions, the `Abort_Signal` exception, an **`at end`** handler should not propagate it to the surrounding execution unless this is done as part of undeferring abortion.

For constructs implemented using matched initialization/finalization TRTS calls like `Enter_Master` and `Complete_Master`, there is a similar problem at the beginning of the construct. `Complete_Master` must be called to leave a master, but the master does not begin until `Enter_Master` is called. If `Enter_Master` is called before entering the block protected by the handler that executes `Complete_Master`, an abortion could occur between the return from `Enter_Master` and the beginning of the block. `Complete_Master` would never be called, resulting in undefined behavior. If `Enter_Master` were protected by the handler, an abortion occurring before `Enter_Master` was called would be handled, resulting in the execution of an unmatched `Complete_Master`.

This problem can be addressed by deferring abortion before entering the block and undeferring it only after the initialization call returns:

```
Defer_Abortion;
declare
  Enter_Master;
  Undefer_Abortion;
  ... mixed master and other object declarations
begin
  ... code for the sequence of statements of the master
at end
  Complete_Master;
end;
```

This is awkward, particularly for procedures, where the `Defer_Abortion` call would have to be made before each procedure call. As a shorthand notation for this technique **`at start`** handler is therefore defined which can be attached to a block, and which will always be executed before the block starts. It defers abortion, executes the specied operations, enters the block, and undefers abortion:

```
at start
  Enter_Master;
declare
  ... mixed task and other object declarations
begin
  ... code for the sequence of statements of the task body
at end
  Complete_Master;
end;
```

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6.5.4 Example of Use of Masters

The following is an example inteded to illustrate how masters enter into the creation of tasks. In this example, the locations of key TRTS calls and implicit variables that must be allocated by the compiler are indicated in comments. For simplicity of exposition, code to assure that all exit and cleanup calls are made in the presense of exceptions is omited.

```
package Pkg_1 is
  task T1;
  -- Create_Task(...Master=> Current_Master,...);
  task type Type_2;
  procedure P;
end Pkg_1;

package body Pkg_1 is
  X: Integer:= 1;
  task body T1 is
  begin
    P;
  --at end
    -- Complete_Task;
  end T1;
  task body Type_2 is
  begin
    if X=1 then X:=X+1; P; end if;
  --at end
    -- Complete_Task;
  end Type_2;
  --at start
    -- Enter_Master;
  procedure P is
    T2: Type_2;
    -- Create_Task(...,Master=> Current_Master,...);
    type Acc_T is access Type_2;
    -- Acc_T_Master: Master_ID:= Current_Master;
    -- is master of tasks created via allocators
    -- returning type Acc_T.
  begin
    --at start
      -- Enter_Master;
    B: declare
      task T3;
      -- Create_Task(..., Master=> Current_Master,...);
      task body T3 is
        procedure R is
          --a trivial master, since it has no local tasks
          --or task access types.
          X: Acc_T;
          --at start
            -- Enter_Master;
          procedure S is
            T4: Type_2;
            -- Create_Task(...,Master=> Current_Master,...);
          begin
            X:= new Type_2;
            -- Create_Task(...,Master=> Acc_T_Master,...);
```

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```
340             --at end
341               -- Complete_Master;
342             end S;
343           begin
344             S;
345           end R;
346         begin
347           R;
348         --at end
349           -- Complete_Task;
350         end T3;
351       begin ...
352       --at end
353         -- Complete_Master;
354       end B;
355     --at end
356       -- Complete_Master;
357     end P;
358   begin ...
359   end Pkg_1;

360   with Pkg_1;
361   --at start
362     -- Enter_Master;
363   procedure Main is
364     T5: Pkg_1.Type_2;
365     -- Create_Task(...,Master=>Current_Master,...);
366   begin null;
367   --at end
368     -- Complete_Master;
369   end Main;
```

370 6.6 Activation Chains

371 The `Activation_Chain` type is a limited private type used to specify sets of tasks
372 that are to be activated together. Activation chain objects are initialized to a state
373 that specifies the empty set of tasks. An object of the type `Activation_Chain` is
374 called an *activation chain*.

375 6.6.1 RTS Responsibilities

376 An activation chain object is *valid for a task* `T` if it is either empty or specifies a set
377 of unactivated tasks, all of which were created by calls to `Create_Task` made by `T`.

378 The TRTS shall define the representation of the `Activation_Chain` type, and en-
379 sure that the default initial value specifies the empty set. If the implementation
380 of activation chains involves dynamic storage allocation (other than on the runtime
381 stack), the the TRTS shall provide appropriate finalization to ensure all storage for
382 activation chain objects is recovered on scope exit.

383 6.6.2 RTS Notes

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The `Activation_Chain` type may be implemented by the TRTS in several different ways. In particular, if tasks are chained together directly, an activation chain may be represented via a reference to the task control block of the last task in the chain.

If separate TRTS storage is used for the chain `Activate_Tasks` can reclaim this storage in the normal case, and either `Expunge_Unactivated_Tasks` or `Complete_Master` can reclaim it in the failure case.

6.6.3 Compiler Responsibilities

The compiler shall pass a valid activation chain to each call to `Create_Task`.

6.6.4 Compiler Notes

The compiler should create an object of type `Activation_Chain` for each master and for each allocator that creates tasks. This chain should be passed to each call to `Create_Task` made by that master or allocator. Either `Activate_Tasks` or `Expunge_Unactivated_Tasks` should be called with this object as an argument before the master's body starts execution or before execution or elaboration proceeds beyond the allocator.

The compiler should only put dependents of a single master on an activation chain. Otherwise, activations might occur in the wrong order.

6.7 Task Creation

The `Create_Task` operation is used to create an unactivated task.

6.7.1 RTS Responsibilities

A call to `Create_Task` shall create a new unactivated task and return a valid task ID representing this task, in the parameter `Created_Task`. The newly created task shall not start execution until it is activated by a call to `Activate_Tasks`. The effects of the parameters shall be as follows:

1. `Size` shall be the initial stack size of the task. The TRTS is responsible for (eventually) allocating space for the stack and initializing it as necessary to activate the task. If `Size` is less than the minimum required to begin activation of a task, the TRTS may round the actual stack size up to that minimum.

2. `Priority` shall be the initial base priority for the task, if it is the range `System.Any_Priority`. If the value is `Unspecified_Priority`, the initial base priority of the task shall be the active priority of the activator, if any; otherwise, `System.Default_Priority`.

   The value `Unspecified_Priority` shall be defined by the TRTS, and shall be outside of the range `System.Any_Priority`.

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3. The `Num_Entries` parameter shall be the number of entries of the task. (How entries are counted is explained further in Section 8.2).

4. `Master` shall be the master on which the task directly depends. The effect is undefined if `Master` is not a valid master ID for the calling task. If `Master` was created by a call made to `Enter_Master` while abortion was deferred, the TRTS may raise `Program_Error` without creating a task.

5. `State` shall be the initial state of the new thread of control. Establishing this state in the new thread of control will be equivalent to calling a procedure with a single argument of type `System.Address` within the context of that thread.

6. `Uninterpreted_Data` is uninterpreted data passed as the actual parameter of the procedure called through `State`.

7. The actual object corresponding to `Elaborated` shall be the *elaboration-check variable* associated with the task. `Activate_Tasks` shall raise `Program_Error` if the value of this object is `False` when `Activate_Tasks` is called (later) for the chain containing this task.

8. The task shall be added to the activation chain specified by the `Chain` parameter. The effect is undefined if the chain already contains a task with a different activator.

The TRTS shall create a new task unless it has insufficient resources to do so. If the TRTS does not have sufficient resources to create the new task it shall raise `Storage_Error`. (Resources needed to create a new task include (at least) storage for the task's runtime stack and storage for TRTS internal data structures.)

6.7.2 RTS Notes

Before returning control to the calling task, the TRTS should initialize all the data structures that it needs in order to be able to handle any of the other TRTS calls that take a `Task_ID` parameter, and save the value of the `State` parameter so that it will become the initial state of the new task when it begins activation.

The TRTS may allocate and initialize the stack any time between the call to `Create_Task` and the activation of the task.

6.7.3 Compiler Responsibilities

The compiler shall provide correct parameters to `Create_Task`. In particular:

1. `Size` shall be a non-negative value.

2. `Priority` shall be `Unspecified_Priority` or a value in the range `Any_Priority`.

3. `Master` shall be a valid master ID for the calling task.

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4. `State` shall permit a procedure with a single `System.Address` parameter to be called with the value of `Uninterpreted_Data`.

5. `Chain` shall be a valid activation chain for the calling task.

6. `Created_Task` shall be discriminated with a value enough to cover the range of entry indices which the compiler will use for rendezvous operations for that task, as defined in Section 8.2.

6.7.4 Compiler Notes

Ada 9X allows tasks to be declared with discriminants. Some means must be provided to allow a task to obtain the values of its discriminants when it is created. For this purpose, `Create_Task` takes a value of type `System.Address` in its `Uninterpreted_Data` parameter and passes it to the compiler-generated procedure which executes the task body.

Exactly how the compiler-generated code uses the address parameter is up to the compiler. All tasks are assumed to share a common virtual address space, so this address can be used as a pointer to a data block containing the discriminants. The compiler can translate the formal discriminant specification of the task to an equivalent record type, with a component for each discriminant, and pass the address of this record as the `Uninterpreted_Data` parameter of `Create_Task`.

The value of the `Priority` parameter should be `Unspecified_Priority` if the initial base priority of the task is not specified by pragma. This allows the TRTS to choose the priority.

The variable passed as `Elaborated` should be initialized to `False` by compiler-generated code when the task type declaration (possibly anonymous) corresponding to the new task is elaborated. The value of this object should be changed to `True` by the compiler-generated code when the task body is elaborated.

Creation of the environment task requires special consideration, and is discussed further below.

6.7.5 Rationale

The `Priority` parameter does not have a constraint, because Ada does not permit raising an exception for incorrect parameters to pragmas, and because it allows use of an out-of-range value to indicate there is no priority specified. The priority of a task where no priority is specified is the priority at which it was created, that is, the priority of the activator at the time that it calls `Create_Task`. This is most easily calculated by the TRTS, so the priority value `Unspecified_Priority` is provided to indicate that no `Priority` pragma was specified for a task.

The `Elaborated` parameter is provided to make implementation of elaboration checks more convenient for the compiler. The compiler cannot always determine statically which tasks will be activated and what their task types are. For example,

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consider the elaboration of a record containing task objects within variant components that depend on the results of function calls; the determination of which tasks are created is made at run time. Without this feature, in those cases where elaboration checks cannot be performed at compile time the compiler would have to generate additional data structures or a special runtime routine to perform these checks for each unconstrained type with task components.

Note also that AI-00149 requires that when activating a list of tasks, if the activation of one task raises `Program_Error` as a result of the task body not being elaborated yet, no task shall be allowed to proceed. Therefore, in order to avoid another synchronization point in the created tasks (after they check their elaboration status), the activator task is required to make the check for all tasks in the chain before beginning the activation of any of them. Furthermore, it is not clear that on all systems one can rely on the task code itself to check its own elaboration status.

6.8 Task Activation

`Activate_Tasks` causes a set of tasks to begin activation.

6.8.1 RTS Responsibilities

`Activate_Tasks` shall cause the TRTS to start activation of the set of tasks specified by the `Chain` parameter. There shall be no return from this call until all tasks in the set complete activating, normally or through an exception, or all are determined by the TRTS not to be activatable, because one or more of the corresponding task bodies has not been elaborated. If all tasks in the set complete activation normally, the call shall return normally and update the `Chain` parameter to designate the empty set. Otherwise, if one or more of the tasks has not completed activating normally, the return shall be via propagation of `Tasking_Error` or `Program_Error`, as specified by the Ada language definition.

Note that `Activate_Tasks` removes all tasks that start activation from the chain parameter, including the case where one or more tasks fails during activation and `Activate_Tasks` returns `Tasking_Error`. Since `Activate_Tasks` would be unable to propagate this exception and still modify the chain parameter directly, a pointer to the `Activation_Chain` object is passed instead. Note also that all tasks on the chain start activation unless one of the task bodies has not been elaborated, in which case none of the tasks begin activation. This means that the chain will either be returned empty or unmodified. a

The effect of `Activate_Tasks` is undefined if `Chain` is not valid for the calling task, or if the chain contains any tasks that depend on a master that has become invalid (due to exit from the master).

6.8.2 Compiler Responsibilities

The activation chain specified by the `Chain` parameter shall be valid for the task that calls `Activate_Tasks`.

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`Chain` shall not contain any tasks that depend on a master which has become invalid.

6.8.3 Compiler Notes

The compiler should create a variable of the `Activation_Chain` type for each non-trivial task master, and should create additional chains to handle the evaluation of allocators that create (and activate) tasks. Whenever it creates a task, the compiler should know the master or allocator evaluation to which it belongs, and pass the appropriate chain. The compiler needs to keep track of which chains need to be activated and generate a call to `Activate_Tasks` for each activation chain at the appropriate location.

6.9 Expunction of Unactivated Tasks

`Expunge_Unactivated_Tasks` expunges a set of unactivated tasks, specified by an activation chain.

6.9.1 RTS Responsibilities

The TRTS shall expunge the tasks on `Chain`. The effect of this call shall be that all the tasks in the specified group become terminated, the associated task IDs become invalid, and all associated resources shall be recovered. The effect is undefined if the activation chain is not valid for the calling task.

6.9.2 Compiler Responsibilities

The activation chain specified by the `Chain` parameter shall be valid for the task that calls `Activate_Tasks`.

`Chain` shall not contain any tasks that depend on a master which has become invalid.

The compiler shall make no further use of task IDs identifying tasks on the activation chain. If any further reference is made to these tasks, the compiler should arrange for a value of `Null_Task` to be substituted for the value originally returned by `Create_Task`.

6.9.3 Compiler Notes

Ada requires a task to be terminated before leaving its master. This includes tasks which were never activated. There must be a call to either `Complete_Master` or `Complete_Task` before exiting the master, and unactivated tasks could be terminated, and any resources allocated to the reclaimed, by this calls. However, tasks can be created by packages and allocators, and these constructs are not masters. Failure of the elaboration of a package spec, or of an allocator, can leave resources allocated to unactivated tasks for as long as the master exists. In the case of library level packages, this is essentially for the lifetime of the program. One possible result is storage leakage.

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Expunge_Unactivated_Tasks allows an implementation to reclaim resources allocated for unactivated tasks upon the failure of a package or allocator that creates tasks. This also allows storage allocated by a failed allocator to reclaim storage immediately, since it must first finalize any subcomponents elaborated before the point of failure, and all of the tasks that it has created must be terminated before finalization can take place.

6.9.4 Example of Task Allocation and Activation

Consider the following declare statement:

```
B:  declare
    task type Task_Type is
      pragma Priority(9);
    end Task_Type;
    type Task_Pointer is access Task_Type;
    for Task_Type'Storage_Size use 200;
    task body Task_Type is ...
    begin ... end Task_Type;
    T: Task_Pointer;
    procedure Inner is
    begin
      T:= new Task_Type;
    end Inner;
          ... other local declarations of B.
  begin
     Inner;
  end B
```

One way to implement this statement is to put **at end** handlers around the allocation to make sure that Expunge_Unallocated_Tasks is always called, as explained in Section 6.5.3. If the call to Activate_Tasks is executed, it will remove all tasks from the chain, in which case calling Expunge_Unactivated_Tasks on it will have no effect.

```
B:  declare
    Task_Type_Elaborated: aliased Boolean:= False;
    Task_Master: Master_ID:= Current_Master;
    procedure Task_Type_Body is ...
    begin ... end Task_Type_Body;
    Task_Type_Elaborated:= True;
    T: Task_ID;
    procedure Inner is
    begin
      declare
        Local_Chain: aliased Activation_Chain;
      begin
         --... code to put initialization data into Init_Data.
         -- For GNAT:
         -- Init_Data:= Task_Type_Body'Access;
         Create_Task(
             Size=> 200,
             Priority=>  9,
             Num_Entries=> 0,
             Master=> Task_Master;
```

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```
615              State=> Init_Data,
616              Uninterpreted_Data=> System.Null_Address,
617              Chain=> Local_Chain,
618              Elaborated=> Task_Type_Elaborated'Access,
619              Created_Task=> T);
620           Activate_Tasks(Local_Chain'Access);                              |a
621        at end
622          Expunge_Unactivated_Tasks(Local_Chain);
623        end;
624      end Inner;
625            ... other local declarations of B.
626    begin
627       Inner;
628    end B
```

Another possibility is to make `Activation_Chain` part of a finalizable object. Since `Local_Chain` is the only declaration in the declare block used to enclose both the creation and activation of the task, no other finalization routines can execute before its finalization[1]. The translation is the same, except that the inner declare block does not need handlers:

```
B:  declare
    Task_Type_Elaborated: aliased Boolean:= False;
    Task_Master: Master_ID:= Current_Master;
    procedure Task_Type_Body is ...
    begin ... end Task_Type_Body;
    Task_Type_Elaborated:= True;
    T: Task_ID;
    procedure Inner is
    begin
      declare
        type Finalizable_Chain is new Controlled with record
          A: aliased Activation_Chain;                                        |a
          end record;
        procedure Finalize(Object: in out Finalizable_Chain) is
        begin
          Expunge_Unactivated_Tasks(Object.A);
        end Finalize;
        Local_Chain: Finalizable_Chain;
      begin
         --... code to put initialization data into Init_Data.
         -- For GNAT:
         -- Init_Data:= Task_Type_Body'Access;
         Create_Task(Size=> 200,
           Priority=>  9,
           Num_Entries=> 0,
           Master=> Task_Master;
           State=> Init_Data,
           Uninterpreted_Data=> System.Null_Address,
           Chain=> Local_Chain.A,
           Elaborated=> Task_Type_Elaborated'Access,
           Created_Task=> T);
         Activate_Tasks(Local_Chain.A'access);                                |a
      end;
```

[1] If there were any such finalization routines, they would be able to see the tasks as unterminated, contrary to Ada semantics.

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```
669     end Inner;
670           ... other local declarations of B.
671   begin
672      Inner;
673   end B
```

674 6.9.5 Example of Activation and Creation

675 For an example of creation and activation through the elaboration of a task object,
676 consider the following block statement, containing two local tasks.

```
677 B: declare
678           ... other local declarations of B.
679      task type TT(D: Integer) is
680        pragma Priority(9);
681      end TT;
682      for TT'Storage_Size use 200;
683      T1: TT(1);
684      task T2 is
685        entry E1(1..10);
686        entry E2;
687        pragma Priority(10);
688      end T2;
689      for T2'Storage_Size use 100;
690      task body T1 is ...
691      begin ... end T1;
692      task body T2 is ...
693      begin ... end T2;
694           ... other local declarations of B.
695    begin ... sequence of statements of B.
696    end B;
```

697 The structure of the translated code for this block statement is given below.

```
698 ... code for elaboration of other local declarations of B
699 T1_ID, T2_ID: Task_ID;
700 T1_Elaborated: constant aliased Boolean:= False;
701 T2_Elaborated: constant aliased Boolean:= False;
702 Local_Chain: aliased Activation_Chain;
703 type Discriminant_Record is record
704   D: Integer;
705   end record;
706 Discriminant_1: Discriminant_Record;
707 procedure T1(Discriminants: System.Address) is ...
708 begin ... end T1;
709 T1_Elaborated:= True;
710 procedure T2(Discriminants: System.Address) is ...
711 begin ... end T2;
712 T2_Elaborated:= True;
713 ...
714 at start
715   Enter_Master;
716 begin
717   ...
718   -- code to put initialization data into INIT_DATA,
719   -- including the starting address of T1.
```

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```
720    -- For GNAT (See Section 1.7.1.1): Init_Data:= T1'Access;

721    Discriminant_1.D:= 1;
722    Create_Task(
723        Size=> 200,
724        Priority=>  9,
725        Num_Entries=> 0,
726        Master=> Current_Master,
727        State=> Init_Data,
728        Uninterpreted_Data=> Discriminant_1'Address,
729        Chain=> Local_Chain,
730        Elaborated=> T1_Elaborated'Access,
731        Created_Task=> T1_ID);

732    -- code to put initialization data into Init_Data.
733    -- For GNAT: Init_Data:= T2'Access;

734    Create_Task(
735        Size=> 100,
736        Priority=> 10,
737        Num_Entries=> 11,
738        Master=> Current_Master,
739        State=> Init_Data,
740        Uninterpreted_Data=> System.Null_Address,
741        Chain=> Local_Chain,
742        Elaborated=> T2_Elaborated'Access,
743        Created_Task=> T2_ID);
744    ... code for elaboration of other local declarations of B
745    Activate_Tasks(Local_Chain'access);
746    ... code for sequence of statements of B.
747 at end
748    Complete_Master;
749 end;
```


750 In this example, `Ti_ID` is an implicit local variable to receive the task ID identifying
751 the new task `Ti`, and `Init_Data` specifies how to call the `Ti` task procedure (e.g. a
752 procedure access value). `Ti` is a procedure used to implement the body of the task
753 `Ti`. As explained in Section 6.5.3, **`at start`** and **`at end`** handlers are used to express
754 the requirement that `Complete_Master` must be called if and only if `Enter_Master`
755 is called, even in the presence of exceptions or abortion.

756 6.10 Activation Completion

757 `Complete_Activation` notifies the TRTS that the calling task has completed acti-
758 vation.

759 6.10.1 RTS Responsibilities

760 If this is the last task to complete activation in a chain of tasks, the TRTS shall release
761 the activator of the chain to resume execution. The effect of calling `Complete_-`
762 `Activation` from a task that has not started activation (i.e. is not on an activation
763 chain for which `Activate_Tasks` has been called), or of calling it more than once
764 from a given task, is undefined.

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765 ### 6.10.2 Compiler Responsibilities

766 The compiler will assure that `Complete_Activation` is called no more than once
767 from each task. It will assure that it is called only from tasks that have started
768 activation.

769 ### 6.10.3 Compiler Notes

770 `Complete_Activation` should be called (exactly once) by each task when it has just
771 completed activation, after the elaboration of its own declarative part is complete.

772 `Complete_Activation` should not be called by a task that fails to complete acti-
773 vation, due to an exception being raised and not handled during the activation. Its
774 functions will be subsumed by `Complete_Task` in this case.

775 ## 6.11 Task Completion

776 `Complete_Task` notifies the TRTS that the body of the calling task has "completed"
777 in the Ada language sense. (The task ordinarily is not completely finished executing,
778 since it may still need to wait for dependents and then execute finalization code.)

779 ### 6.11.1 RTS Responsibilities

780 If the task has dependents, the actions of the TRTS for `Complete_Task` shall include
781 all actions required for `Complete_Master`. In addition, the TRTS shall perform all
782 actions required by the Ada language for completion of a task.

783 1. Check for pending calls on entries of the task, and raise `Tasking_Error` in any
784 such calling tasks.

785 2. Arrange to raise `Tasking_Error` in the activator, if the present task is not
786 finished activating (by calling `Complete_Activation`).

787 3. Set the `Callable` attribute of the task to `False`.

788 The effect of `Complete_Task` is undefined if there is a rendezvous that was begun by
789 a call to `Selective_Wait` or `Accept_Simple` but has not yet been completed by a
790 matching call to `Complete_Rendezvous` or `Exceptional_Complete_Rendezvous`.
791 Its effect is also undefined if called after a call to `Lock` or `Lock_Read_Only` without
792 a corresponding call to `Unlock`. Its effect is also undefined if it is called between
793 a call to `Task_Entry_Call` and the corresponding `Cancel_Task_Entry_Call`, or
794 `Protected_Entry_Call` and the corresponding `Cancel_Protected_Entry_Call`. |a

795 ### 6.11.2 RTS Notes

796 `Complete_Task` has no parameters. All the context information needed to process
797 this call needs to be kept by the TRTS.

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Ada requires a task to be completed before finalization. Since finalization will need access to objects declared in the task, `Complete_Task` cannot in general reclaim the resources used for such objects. This usually means that `Complete_Task` will be unable to deallocate the task's stack. This must be deferred until task termination.

6.11.3 Compiler Responsibilities

The compiler shall ensure that `Complete_Task` is called (once) after the task body completes, whether normally, via a **`return`**, **`exit`** or **`goto`** statement, or through propagation of an exception, but prior to finalization of the task. It shall assure that all nesting runtime calls are balanced on all control paths leading to the call to `Complete_Task`; for example, if the task has made a call to `Lock`, there must be a corresponding call to `Unlock` before calling `Complete_Task`.

6.11.4 Compiler Notes

As with `Complete_Master`, the requirement that `Complete_Task` be called before an exception is propagated out of a task body cannot be expressed directly in Ada (see Section 6.5.3). Examples of task completion are shown using Ada with extensions.

`Complete_Task` will wait for completion of all tasks dependent on it, even those created by masters executed by the task. However, it will not ensure the finalization of such internal masters, and those masters must wait on their own dependent tasks before finalization. Therefore, `Complete_Task` should not be called until all of the masters it has started executing have been completed and left.

6.12 Task Termination

The TRTS "terminates" a task, in the Ada language sense, when the task body is left. At this point, the task has completed execution, all of its dependents are terminated, and all of its finalization has been performed.

6.12.1 RTS Responsibilities

The TRTS shall set the `Terminated` attribute of the task to `True` on leaving the body of a task.

The effect of leaving a task body is undefined if the task has not previously called `Complete_Task`. Its effect is also undefined if finalization has not been performed on the task.

6.12.2 RTS Notes

The TRTS calls the procedure representing the task body to elaborate the task body declarations and execute its sequence of statements. This procedure returns to the point of call in the TRTS when the task completes, after calling `Complete_Task`, as specified above, and after all finalization has been completed. The TRTS terminates the task at this point.

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Since finalization must be complete when the body is left, there should be no further reference to objects declared in the task body. The TRTS can therefore deallocate these objects, usually by deallocating the task stack.

6.12.3 Compiler Responsibilities

The compiler shall implement the body of a task as a procedure (see Section 1.7.1).

6.12.4 Example of Task Activation, Completion, and Termination

Consider a task body such as:

```
task body T is
  -- declarations
begin
  -- sequence of statements
end T;
```

`Complete_Activation`, `Complete_Task`, and `Terminate_Task` can be used in a task body procedure to control the lifetime of this task as follows:

```
procedure T_Body(Discriminants: System.Address) is
begin
  ... code for elaboration of the task's local declarations
  Complete_Activation;
  ... code for the sequence of statements of the task body.
at end
  Complete_Task(T);
end T_Body;
```

6.13 Terminated Attribute

The `Terminated` function shall return the value of the language-defined `Terminated` task attribute

6.13.1 RTS Responsibilities

`Terminated` shall return `True` if its argument represents a task that has been terminated, or is `Null_Task`; otherwise, it returns false. The effect of calling `Terminated` with an invalid task ID that is not `Null_Task` is undefined.

6.13.2 Compiler Responsibilities

The compiler shall assure that the argument to `Terminated` is a valid task ID or `Null_Task`.

6.13.3 Compiler Notes

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For each instance of the task attribute `T'Terminated`, the compiler should generate a call to the `Terminated` function with a task ID identifying `T` as its argument. If the compiler-generated code has expunged a task before its master is left, this task is also terminated, though the task ID assigned to it is no longer valid. In such cases, the compiler-generated code can translate all references to the task to `Null_Task`, or simply replace its `Terminated` attribute with `True`.

6.14 Environment Tasks

The environment task is implicitly created by the elaboration of the Ada TRTS packages.

6.14.1 RTS Responsibilities

The TRTS shall treat the thread of control that elaborates the TRTS packages as the environment task. After the elaboration of the TRTS packages is complete, it shall be allowable for this thread of control to make any legal TRTS call.

6.14.2 RTS Notes

When an environment task completes, it should wait for the termination of tasks created by library packages, which are dependent on the environment task. The TRTS should also perform whatever additional cleanup is needed to shut down the partition. This might include closing of any open files. A call-out for user-defined finalization should perhaps be provided, to execute after all the Ada 9X object finalization.

Immediate shut-down of an entire process, as via the POSIX *exit()* operation, should bypass normal finalization and termination processing. Otherwise, there is risk of a program never being able to shut down at all.

6.14.3 Compiler Responsibilities

The elaboration order of the load-image must ensure that the TRTS code is elaborated to completion, by a single thread of control, before the execution of any code that makes TRTS calls.

6.14.4 Compiler Notes

Each nontrivial master is required to be bracketed by calls to `Enter_Master` and `Complete_Master`. This includes the main program procedure, which is executed by the environment task. Tasks created by the main program procedure are dependent on it, and not (directly) on the environment task, and the main program procedure must wait for them to terminate before returning to the TRTS.

The compiler should translate an Ada program (or active partition) into a program that can be loaded as a single-threaded process under the host operating system.

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Section 7: Delays

This section describes the TRTS support for various forms of **delay** statements.

7.1 Package Specifications

```
package System.Real_Time.Delays is

   protected Delay_Until_Object is
     entry Wait(T: Time);
   end Delay_Until_Object;

end System.Real_Time.Delays;

package Calendar.Delays is

   protected Delay_Object is
     entry Wait(T: Duration);
   end Delay_Object;

   protected Delay_Until_Object is
     entry Wait(T: Time);
   end Delay_Until_Object;

end Calendar.Delays;
```

7.1.1 RTS Responsibilities

Expressions of *time types* can be used in **delay until** clauses. These must be non-limited types on which an ordering function "<" is defined. These restrictions are needed for the translation of some types of **delay** clauses, e.g. those described in Section 8.7.5.

A package *time_package* provided by the TRTS which defines a time type shall have a child package *time_package*.Delays which defines Delay_Until_Object with a Wait entry which takes that time type as a parameter. A call to the Wait entry will be queued for the amount of time specified by its argument, as measured by the Clock function in the corresponding package. After this time has arrived, the entry call will be dequeued. The Wait entry body has no other effect.

Relative delays are supported by the Delay_Object defined in System.Real_Time package. The GNARLI provides a similar Calendar.Delays.Delay_Object which defines a Wait entry that takes a Duration parameter.

7.1.2 Compiler Notes

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The compiler should generate a simple protected entry call on `Calendar.Delays.Delay_Object.Wait` for a **delay** statement, passing its `Duration` expression to the entry.

The compiler should generate a simple protected entry call on *time_package*`.Delays.Delay_Object.Wait` for a **delay until** statement, where *time_package* is the Ada package defining the time type of the delay expression.

Asynchronous delay statements should be implemented as asynchronous protected entry calls to the appropriate `Wait` entry.

Timed entry calls should be implemented as asynchronous entry calls, where the abortable part is a simple **delay** or **delay until** statement.

Select statements should be implemented as asynchronous protected entry calls to the appropriate `Wait` entry, with the `Selective_Wait` TRTS call as the abortable part.

7.1.3 Example of a Simple Relative Delay Statement

Consider a simple relative delay statement:

```
delay T;
```

This can be implemented by the following code fragment:

```
Calendar.Delays.Delay_Object.Wait(T);
```

7.1.4 Example of a Simple Absolute Delay Statement

Consider a simple absolute delay statement on a time type implemented by the package `Some_Time_Type`.

```
declare
  T: Some_Time_Type.Time:= Some_Time_Type.Time_Of(1993, 12, 10);
begin
  delay until T;
end;
```

Note that there need not be a `Time_Of` function in this package; it is only shown to indicate some method of obtaining a time type value. This can be implemented by the following code fragment:

```
declare
  T: Some_Time_Type.Time:= Some_Time_Type.Time_Of(1993, 12, 10);
begin
  Some_Time_Type.Delays.Delay_Until_Object.Wait(T);
end;
```

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7.1.5 Example of an Asynchronous Delay Statement

Consider the following asynchronous **select** statement with a **delay until** trigger:

```
declare
  T: Some_Time_Type.Time:= Some_Time_Type.Time_Of(1993, 12, 10);
begin
  select
    delay until T;
  then abort
    S;
  end;
end;
```

This can be implemented by the following code fragment:

```
declare
  T: Some_Time_Type.Time:= Some_Time_Type.Time_Of(1993, 12, 10);
begin
  select
    Some_Time_Type.Delays.Delay_Until_Object.Wait(T);
  then abort
    S;
  end select;
end;
```

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Section 8: Rendezvous

The TRTS operations required to implement rendezvous are provided by the `System.Rendezvous` package.

8.1 Package Specification

```
with System.Compiler_Exceptions;
with System.Task_IDs; use System.Task_IDs;
with System.Protected_Objects;
package System.Rendezvous is

   Null_Task_Entry: constant:= 0;
   Max_Task_Entry: constant:= ...; -- value defined by TRTS.
   type Task_Entry_Index is range Null_Task_Entry..Max_Task_Entry;

   procedure Task_Entry_Call(
       Acceptor: Task_ID;
       E: Task_Entry_Index;
       Uninterpreted_Data: System.Address;
       Mode: System.Protected_Objects.Call_Modes;
       Rendezvous_Successful: out Boolean);
   procedure Call_Simple(
       Acceptor: Task_ID;
       E: Task_Entry_Index;
       Uninterpreted_Data: System.Address);
   procedure Cancel_Task_Entry_Call(Cancelled: out Boolean);
   procedure Requeue_Task_Entry(
       Acceptor: Task_ID;
       E: Task_Entry_Index;
       With_Abort: Boolean);
   No_Rendezvous: constant:= 0;
   type Select_Modes is (Simple_Mode, Else_Mode, Terminate_Mode);
   Max_Select: constant:= implementation-defined;
   type Select_Index is range No_Rendezvous..Max_Select;
   type Select_Index_Access is access all Select_Index;                    a
   subtype Positive_Select_Index is
       Select_Index range 1 .. Select_Index'Last;
   type Accept_Alternative is record -- should be packed
      Null_Body: Boolean;
      E: Task_Entry_Index;
      end record;
   type Accept_List is
       array(Positive_Select_Index range <>) of Accept_Alternative;
   -- Lower index bound shall be 1.
   procedure Selective_Wait(
       Open_Accepts: Accept_List;
       Select_Mode: Select_Modes;
       Uninterpreted_Data: out System.Address;
```

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```
45          Index: Select_Index_Access);                                    |a
46      procedure Accept_Call(
47          E: Task_Entry_Index;
48          Uninterpreted_Data: out System.Address);
49      procedure Accept_Trivial(E: Task_Entry_Index);
50      function  Task_Count(E: Task_Entry_Index) return Natural;           |a
51      function  Callable(T: Task_ID) return Boolean;
52      procedure Complete_Rendezvous;
53      procedure Exceptional_Complete_Rendezvous(
54          Ex: Compiler_Exceptions.Exception_ID);
55  end System.Rendezvous;
```

56 The description of this package is divided into several subsections.

57 8.2 Entry Indices

58 8.2.1 RTS Responsibilities

59 The `Task_Entry_Index` type is used to identify specific entries of a task. An entry
60 index value `i` is *valid for a task* `T` if `i` is in the range `1..Num_Entries`, where `Num_-`
61 `Entries` is the number of entries specified in the call to `Create_Task` that created
62 `T`.

63 8.2.2 Compiler Responsibilities

64 The compiler shall assign a unique positive value of the type `Task_Entry_Index` to
65 each entry of a task type. In assigning indices to entries, each component of an entry
66 family shall be considered to be an entry.

67 The largest entry index value used for a task shall be communicated by the compiler-
68 generated code to the TRTS at the time of task creation, via the `Num_Entries`
69 parameter of `Create_Task`.

70 8.2.3 Compiler Notes

71 It is recommended that the entry indices used for a task be assigned in a compact
72 range, in the order in which the entry declarations appear. The compiler should
73 not leave gaps in the range of entry indices assigned to a task, since the TRTS may
74 allocate storage for a separate entry queue for each entry index value in the range
75 `1..Num_Entries`.

76 The compiler should allocate indices for members of entry families in a fashion
77 similar to the allocation of storage for arrays, as if an entry family were an array of
78 entries, and the entry indices were addresses of individual elements.

79 8.3 Entry Calls

80 The `Task_Entry_Call` operation is used to implement Ada entry calls.

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8.3.1 RTS Responsibilities

This operation shall make an entry call to the task entry specified by the `Acceptor` and `E` parameters. The effect of the operation is undefined unless `Acceptor` is a valid task ID, and `E` is a valid entry index for that task. The value of `Uninterpreted_Data` shall be passed to the acceptor, as the `Uninterpreted_Data` of the corresponding accept operation (see Section 1.9). If `Acceptor` represents a terminated task, this call will raise `Tasking_Error` in the caller.

`Task_Entry_Call` makes different kinds of entry calls according to the value of the `Mode` parameter. The specific actions of `Task_Entry_Call` for the various values of `Mode` are:

1. `Simple_Call` shall make an unconditional entry call.
2. `Conditional_Call` shall make a conditional entry call. It shall set the parameter `Rendezvous_Successful` to `False` if there is no rendezvous, and `True` if there is a successful rendezvous.
3. `Asynchronous_Call` shall make an asynchronous entry call, during which the calling task may continue execution of an abortable sequence of statements.

The detailed TRTS actions required for these calls are:

1. Check that the acceptor is callable, and raise `Tasking_Error` if it is not.
2. For `Conditional_Call`, return immediately if the called task is not ready to accept.
3. For `Simple_Call`, if the called task is not ready to accept, block the calling task.
4. For `Asynchronous_Call`, if abortion is deferred, the TRTS may raise `Program_Error`. Otherwise, if the called task is not ready to accept, queue the calling task on the specified entry.
5. When the called task is ready to accept, transfer the parameter address and begin the rendezvous. (This involves unblocking and possibly adjusting the dispatching priority of the accepting task.) Except for `Asynchronous_Call`, block the calling task.

8.3.2 Compiler Responsibilities

The compiler is required to pass valid parameter values:

1. `Acceptor` must be a valid task ID.
2. `E` must be an entry index in the range `1..Num_Entries`, where `Num_Entries` is the parameter passed to the call to `Create_Task` that created the `Acceptor` task.

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8.3.3 Compiler Notes

With Ada 9X we can have a task continue execution while an entry call is pending. We must take care to put the entry parameters some place where they will not be clobbered. This will probably be done by using the runtime stack. The effect will be similar to treating the asynchronous select statement as a declare-block, where the entry parameter block is a local variable of the declare-block.

`Task_Entry_Call` is intended to be used to implement simple, conditional, or asynchronous entry calls when used with a mode of `Simple_Call`, `Conditional_Call`, or `Asynchronous_Call`, respectively. It can also be used in conjunction with **`delay`** to implement timed entry calls in the `Asynchronous_Call` mode.

The following examples are intended to show how these operations could be used by a compiler to implement conditional entry calls. Examples of the implementation of timed and asynchronous calls are given in Section 8.5.

8.3.4 Example of a Conditional Entry Call

An example of a conditional entry call is:

```
select T.E(X); ... sequence of statements (1).
else           ... sequence of statements (2).
end select;
```

Suppose `E` is the entry of task `A` associated with index 1, and the task ID identifying `T` is stored in the variable `T`. The statement above would then be translated into something like the following, where `S` is a local temporary variable supplied by the compiler:

```
... code to copy the parameter-passing information into location P.
Task_Entry_Call(
    Acceptor=> T,
    E=> 1,
    Uninterpreted_Data=> P,
    Mode=> Conditional_Call,
    Rendezvous_Successful=> S);
if S then ... code for sequence of statements (1).
else            ... code for sequence of statements (2).
end if;
```

8.4 Simple Entry Calls

The `Call_Simple` operation is provided to implement the simple call statement with potentially less overhead than if then `Task_Entry_Call` were used.

8.4.1 RTS Responsibilities

`Call_Simple(T, E, P)` shall be equivalent to

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```
154   Task_Entry_Call(
155       Acceptor=> T,
156       E=> E,
157       Uninterpreted_Data=> P,
158       Mode=> Simple_Call,
159       Rendezvous_Successful=> S)
```

160 where S is a dummy variable whose value is ignored.

161 8.5 Completing Asynchronous Task Entry Calls

162 Cancel_Task_Entry_Call is used to remove a pending asynchronous task entry call
163 from the corresponding queue.

164 8.5.1 RTS Responsibilities

165 Cancel_Task_Entry_Call shall remove the task entry call (that has not already
166 been removed) made by the most recent call to Task_Entry_Call executed by the
167 calling task. If the call has not yet completed, removing it entails cancellation. If
168 the call is not either complete or queued non-abortably, Cancel_Task_Entry_Call
169 shall block until the it is. If the call is actually cancelled, the Cancelled parameter
170 will be True on return; otherwise, it will be False.

171 The effect of calling Cancel_Task_Entry_Call when the last entry call made by
172 the task which has not yet been removed is not an asynchronous task entry call is
173 undefined.

174 8.5.2 Compiler Responsibilities

175 The compiler shall assure that Cancel_Task_Entry_Call is called only when the
176 most recent entry call made by the calling task is an asynchronous protected entry
177 call, and that Cancel_Task_Entry_Call is called only once for this call.

178 8.5.3 Compiler Notes

179 The compiler should call Cancel_Task_Entry_Call exactly once for each asyn-
180 chronous task entry call when the abortable part completes. This applies to the
181 case where the abortable part is aborted by the completion of the triggering task
182 entry call. This implements the semantics of the completion of an abortable part,
183 cancelling the call if queued abortably and waiting for it to complete or be queued
184 abortably if not. In either case, this call removes the current asynchronous task
185 entry call and makes the previous asynchronous entry call (if any) current. This is
186 the reason that this call must be made exactly once for each asynchronous task entry
187 call; it must be called once to remove the call, but calling it again would remove the
188 previous call prematurely, or result in undefined behavior if no such call exists.

189 The compiler should arrange for calls to Task_Entry_Call with an Asynchronous_-

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Call mode parameter and Cancel_Task_Entry_Call to be executed in pairs, even in the presence of asynchronous abortion. This is similar to the requirement for paired execution of Enter_Master and Complete_Master, as explained in Section 6.5.3, and the same extensions to Ada used to express that requirement are used in the following examples of the use of Cancel_Task_Entry_Call

8.5.4 Example of an Asynchronous Task Entry Call

Consider the following asynchronous **select** statement on a task entry:

```
select
  T.E;
  S1;
then abort
  S2;
end select;
```

This call might be translated as:

```
at start
  declare
    Dummy: Boolean;
  begin
    Task_Entry_Call(
        Acceptor=> T,
        E=> E,
        Uninterpreted_Data=> System.Null_Address;
        Mode=> Asynchronous_Call,
        Rendezvous_Successful=> Dummy);
  end;
declare
  Cancelled: Boolean;
begin
  S2;
at end
  Cancel_Task_Entry_Call(Cancelled);
end;
if not Cancelled then
  S1;
end if;
```

Task_Entry_Call either starts or queues the entry call abortably and returns, at which time the abortable part S2 is executed. If the abortable part completes normally, it attempts to cancel the entry call using Cancel_Task_Entry_Call. If the entry call is queued abortably, it will be canceled; otherwise, Cancel_Task_Entry_Call waits until the call is complete. If the entry call completes successfully, the task will be marked with the nesting level of this asynchronous select and an Abort_Signal pseudo-exception will be raised by Complete_Rendezvous. This will be caught by the compiler-generated exception handler which will call Cancel_Task_Entry_Call.

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8.5.5 Example of a Timed Task Entry Call

The asynchronous form of `Task_Entry_Call` can be used to implement a timed task entry call. For example, the time entry call:

```
select
  T.E;
  S1;
or delay until Calendar.Time_Of(1993, 12, 10);
  S2;
end select;
```

might be translated as:

```
at start
  Task_Entry_Call(
      Acceptor=> T,
      E=> E,
      Uninterpreted_Data=> System.Null_Address,
      Mode=> Asynchronous_Call,
      Rendezvous_Successful=> Dummy);
declare
  Deadline: Some_Time_Type.Time:= Some_Time_Type.Time_Of(1993, 12, 10);
  Timed_Out: Boolean:= False;
  Ex: Exception_ID:= Null_Exception;
begin
  delay until Deadline;
at end
  Cancel_Task_Entry_Call(Timed_Out);
end;
if Timed_Out then
  S2;
else
  S1;
end if;
```

Note that this translation uses a **`delay until`** statement, which itself should be implemented in terms of TRTS operations (see Section 7).

8.6 Requeuing an Entry Call

8.6.1 RTS Responsibilities

`Requeue_Task_Entry` shall requeue the entry call that is currently being served to a specified task entry. The effect is undefined unless the call occurs while the calling task is serving an entry call. The entry call being serviced can be a task or protected entry call (see Section 9).

The effects of the parameters are as follows:

1. `Acceptor` shall be the task on whose entry on which the call is requeued. The effect is undefined unless this is a valid task ID.

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2. `E` shall be the index of the entry on which the call is to be requeued. The effect is undefined unless this is a valid entry index for the specified task.

3. `With_Abort` specifies that the requeued entry call shall be abortable. The call cannot be cancelled until it is queued abortably.

8.6.2 Compiler Responsibilities

`Requeue_Task_Entry` shall only be called by a task that has executed `Selective_Wait` or `Accept_Simple` and before it has executed the corresponding `Complete_Rendezvous` or `Exceptional_Complete_Rendezvous`; or which has executed `Next_Entry_Call` and before it has executed the corresponding `Complete_Entry_Body` or `Exceptional_Complete_Entry_Body`. The compiler shall provide correct parameters; in particular:

1. `Acceptor` shall be a valid task ID.

2. `E` shall be valid for the task corresponding to `Acceptor`.

8.6.3 Compiler Notes

`Requeue_Task_Entry` is intended to implement the Ada **`requeue`** statement to a task entry when executed within an **`accept`** or **`entry`** body. It should be called with `With_Abort` equal to `True` to implement an abortable **`requeue`** statement, and equal to `False` to implement a non-abortable **`requeue`** statement.

8.7 Selective Waits

`Selective_Wait` is provided to implement selective wait statements.

8.7.1 RTS Responsibilities

The effect of the `Selective_Wait` call shall be to implement the form of selective wait statement specified by the `Select_Mode` parameter. The specific effects of the parameters are:

1. `Open_Accepts` shall specify the set of open accept alternatives, on which the calling task is prepared to accept calls. Each component `Open_Accepts(I)` specifies one accept alternative. If `Open_Accepts(I).E` is equal to `Null_Task_Entry`, the accept alternative shall be considered closed. Otherwise, the alternative shall be considered open and ready to accept a call on entry of the calling task with index `Open_Accepts(I).E`. If `Open_Accepts(I).E` is not a valid entry index for the calling task, the effect of the call is undefined. The TRTS must allow the same entry index to appear in more than one component of the list.

   If more than one open accept alternative has a pending entry call, which is chosen depends on the queuing policy implemented by the TRTS. Under FIFO

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queuing, the TRTS will choose one arbitrarily. With priority queuing, more than one accept alternative has a pending entry call, and there is a tie for highest priority, the tie shall be broken in favor of the earlier alternative in the list.

2. `Uninterpreted_Data` is set to the value provided by the caller in its `Uninterpreted_Data` argument, if a rendezvous is successful.

3. The object referenced by `Index` is set to the index of the alternative selected, in the list `Open_Accepts`, if a rendezvous is successful. Otherwise, it is set to `No_Rendezvous`. This parameter is implemented as a pointer to allow it to be updated correctly even if `Selective_Wait` is interrupted by ATC.

The actions performed by the TRTS for the `Selective_Wait` operation are:

1. If a call cannot be accepted immediately:

   (a) For `Else_Mode`, return.

   (b) For `Simple_Mode`, block the accepting task and record that it is waiting for a call to one of the open entries.

   (c) For `Terminate_Mode`, check for possible termination of all the dependent tasks of a master. If the task is not ready for termination, block the accepting task and record that it is waiting for either termination or a call to one of the open entries.

2. If a call can be accepted, raise the acceptor's dispatching priority if necessary, and return control to the acceptor, with the value of `Index` set to indicate which **`accept`** alternative is selected, and the value of `Uninterpreted_Data` set to the value provided by the corresponding entry call operation.

If the call to `Selective_Wait` returns without a rendezvous, the value referenced by `Index` shall be set to `No_Rendezvous`. Otherwise, this value shall be the index of the accept alternative chosen for rendezvous.

The processing of a rendezvous depends on the value of the `Null_Body` component of the accept alternative selected:

1. If `Null_Body` equals `True`, the rendezvous shall be assumed to be *trivial*. If the entry call accepted was initiated by an asynchronous **`select`** statement, the corresponding abortable part shall be aborted before return from `Selective_Wait` (see Section 5); otherwise, the task that made the entry call shall be made eligible for execution before return from `Selective_Wait`.

2. If `Null_Body` equals `False`, the rendezvous is not trivial, and the caller will continue to execute its abortable part (if any) and then block until the acceptor task signals the end of rendezvous, by calling `Complete_Rendezvous` or `Exceptional_Complete_Rendezvous`.

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Between the time that `Selective_Wait` returns with a non-trivial rendezvous selected and the time that the corresponding `Complete_Rendezvous` or `Exceptional_Complete_Rendezvous` is executed, a task is said to be *serving a task entry call*.

8.7.2 Compiler Responsibilities

The lower bound of the array passed for `Open_Accepts` shall be one. The `E` component of every array component must be either `Null_Task_Entry` or a valid entry index for the calling task.

The compiler shall not call `Complete_Rendezvous` or `Exceptional_Complete_Rendezvous` unless it is serving a task entry call.

8.7.3 Compiler Notes

The `Selective_Wait` operation can be called to implement all the forms of selective wait statement. The `Select_Modes` type is used to specify a particular form of selective wait statement. The meanings of the values are:

1. `Simple_Mode`. There are only **`accept`** alternatives.
2. `Else_Mode`. There are only **`accept`** alternatives and an **`else`** clause.
3. `Terminate_Mode`. There are only **`accept`** alternatives and a **`terminate`** alternative.

Ordinarily, the compiler should use the mode that corresponds to the syntactic form of the statement. However, if it can determine that certain guards are closed, it may use a more restrictive form. In particular, the compiler can use `Simple_Mode` for various syntactic forms, including simple **`accept`** statements and selects with **`terminate`** alternatives if it can determine that guards on all the non-**`accept`** alternatives are closed.

Note that the components of the `Open_Accepts` parameter naturally correspond to the alternatives of a selective wait statement, and can be generated in the order in which they appear within the text of the Ada source program. Thus, the compiler will be able to translate a selective wait as a **`case`** statement, using the value returned in the `Index` parameter. Components corresponding to **`accept`** alternatives that have no sequence of statements should have the `Null_Body` field set to `True`. If such an alternative is selected, a trivial rendezvous will be performed before `Selective_Wait` returns.

After a return from `Selective_Wait` for which a nontrivial rendezvous has started, the compiler-generated code is responsible for (later) calling `Complete_Rendezvous` or `Exceptional_Complete_Rendezvous` to end the rendezvous. The value of `Index` and prior knowledge of the value of the corresponding `Null_Body` component of `Open_Selects` can be used to determine whether a nontrivial rendezvous has started.

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The compiler may use the value `Null_Task_Entry` to turn off those alternatives whose guards evaluate to `False` or which are not **accept** alternatives, so as to preserve a one-to-one correspondence between cases and positions in the open entry list. Alternatively, the compiler may simply omit components for non-**accept** entries and **accept** entries that are closed.

The textual order of the **select** alternatives should be reflected in the `Select_Index` field of the corresponding `Open_Accepts` component. This allows the TRTS to implement queuing policies where textual order is significant.

The compiler-generated code for a selective wait should do several things before calling `Selective_Wait`:

1. It should evaluate any guards on the alternatives of the selective wait statement, then if there are entry family members in open alternatives it should evaluate the entry index expressions.
2. If there are no open alternatives, and there is no **else** part, it should raise `Tasking_Error`.
3. If there are no open alternatives, and there is an **else** part, it should transfer control to the **else** part.
4. For each (open) **accept** alternative, it should add a component to `Open_Accepts`.

> **This requires the compiler to do some complicated stuff in code generation. An alternative is to have the TRTS do some of this processing. The reason MRTSI/CIFO chose to put the burden on the compiler was a hope that sometimes this processing could be optimized, using local context.**

In cases where there are no guards, the work can mostly be done at compile time, and `Open_Accepts` can be a constant. This optimization may significantly reduce the execution time of a **select** statement.

8.7.4 Example of a Selective Wait Statement

An example of a selective wait statement and rendezvous code is:

```
select
  when B1=> accept A do S1; end A; T1;
  or when B2=> accept B; T2;
  or when B3=> accept C do S3; end C; T3;
  or when B4=> accept A; T4;
  else T5;
end select;
```

One of several equivalent ways to structure the TRTS calls for this statement would be as follows:

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```
423 declare
424   Q: Accept_List:=
425       ((Null_Body=>False,E=>A),
426       (Null_Body=>True, E=>B),
427       (Null_Body=>False,E=>C),
428       (Null_Body=>True, E=>A));
429   X: aliased Select_Index;                                          a
430   P: System.Address;
431 begin
432   if not B1 then Q(1).E:= Null_Task_Entry; end if;
433   if not B2 then Q(2).E:= Null_Task_Entry; end if;
434   if not B3 then Q(3).E:= Null_Task_Entry; end if;
435   if not B4 then Q(4).E:= Null_Task_Entry; end if;
436   -- or some "optimized" version of the above.

437   Defer_Abortion;                                                   |
438   Selective_Wait(                                                   |
439       Open_Accepts=>Q,                                              |
440       Select_Mode=> Else_Mode,                                      |
441       Uninterpreted_Data=>P, Index=>X'Access);                      |
442   case X is                                                         |
443   when No_Rendezvous=>                                              |
444     Undefer_Abortion;                                               |
445     T5;                                                             |
446   when 1=>                                                          |
447     begin                                                           |
448       Undefer_Abortion;                                             |
449       S1;                                                           |
450     at end                                                          |
451       Exceptional_Complete_Rendezvous(Current_Rendezvous);          |
452     end;                                                            |
453     T1;                                                             |
454   when 2=>                                                          |
455     Undefer_Abortion;                                               |
456     T2;                                                             |
457   when 3=>                                                          |
458     begin                                                           |
459       Undefer_Abortion;                                             |
460       S3;                                                           |
461     at end                                                          |
462       Exceptional_Complete_Rendezvous(Current_Rendezvous);          |
463     end;                                                            |
464     T3;                                                             |
465   when 4=>                                                          |
466     Undefer_Abortion;                                               |
467     T4;                                                             |
468   when others=> raise Program_Error;                                |
469   end case;                                                         a

470 end
```

471 In this example, `Q` is first initialized to a constant vector, then those entries corre-
472 sponding to closed alternatives are "turned off". The creation of `Q` could be optimized
473 by eliminating items corresponding to closed alternatives from `Q`, or if `B1`, `B2`, `B3`,
474 and `B4` are static or loop-invariant, the construction of `Q` could be performed earlier.
475 `P` is used by the translated images of `S1..S4` to access the actual entry parameters.

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The identifier `Current_Exception` stands for the current exception, which the compiler-generated code is responsible for knowing, possibly with the help of the exception-handling portion of the RTS. (Recall that the interface between the exception handler code and the exception handling portion of the RTS is not part of the TRTS interface.) This exception will be raised by the RTS in the calling task at the point of the call. If the **at end** handler was not entered as the result of an exception, `Current_Exception` should be `Null_Exception`, which will not raise an exception in the calling task. Alternately, `Complete_Rendezvous` could be called in cases where no exception has been raised. If the select statement is aborted, `Current_Exception` should be `Abort_Signal`; the TRTS will raise `Tasking_Error` in the caller in this case.

Note that there is no rendezvous code, and no call to `End_Rendezvous` or `Complete_Rendezvous`, for the cases where `Null_Body` is true.

8.7.5 Example of a Selective Wait Statement With Delay Alternatives

`Selective_Wait` can also be used in conjunction with delay protected objects to implement **select** statements with **delay** alternatives. This can be implemented using the `Selective_Wait` operation, with the timeout provided by an asynchronous **select** statement with a `Delay_Object` trigger.

`Selective_Wait` should first be called with `Else_Mode`; otherwise, a very short timeout could cancel the `Selective_Wait` before it could accept a call that was already queued, contrary to Ada semantics. If there is more than one open **delay** alternative, the compiler should choose the alternative with the minimum delay, and use the corresponding delay expression as the argument to the `Wait` entry call used to implement the timeout. If there is more than one minimum delay, the alternative should be chosen according to the queuing policy implemented by the TRTS. If the minimum delay value is negative or zero, there is no need for a timeout, and the `Selective_Wait` operation with mode `Else_Mode` is sufficient.

For example, consider the following selective wait statement.

```
select
  when B1=> accept A do S1; end A; T1;
  or when B2=> accept B; T2;
  or delay until Time1; T3;
  or delay until Time2; T4;
  else T5;
end select;
```

One possible translation of this statement is shown below.

```
declare
  X: aliased Select_Index:= No_Rendezvous;
  Q: Accept_List;
  P: System.Address;
  T: Some_Time_Type.Time;
begin
  Q:= ((Null_Body=>False,E=>A),
       (Null_Body=>True, E=>B));
```


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```
520   if not B1 then Q(1):= Null_Task_Entry; end if;
521   if not B2 then Q(2):= Null_Task_Entry; end if;
522   if Time1 < Time2 then
523     T:= Time1;
524   else
525     T:= Time2;
526   end if;
527   -- or some "optimized" version of the above.

528   -- Try for immediate rendezvous first.

529   at start
530     Selective_Wait(Q, Else_Mode, P, X'Access);                    |a
531   begin
532     case X is
533     when No_Rendezvous=>
534       null;
535     when 1=>
536       S1;
537     when 2=>
538       T2;
539     when others=>
540       null;
541     end case;
542   at end
543     if X /= No_Rendezvous and then not Q(X).Null_Body then
544       Exceptional_Complete_Rendezvous(Current_Rendezvous);
545     end if;
546   end;

547   -- If immediate rendezvous does not work, try waiting for the
548   -- specified delay.

549   if X = No_Rendezvous then
550     begin
551       select
552         delay until T;
553       then abort
554         Selective_Wait(Q, Simple_Mode, P, X'Access);             |a
555       end select;

556       case X is
557       when No_Rendezvous=>
558         null;
559       when 1=>
560         S1;
561       when 2=>
562         T2;
563       when others=>
564         null;
565       end case;
566     at end
567       if X /= No_Rendezvous and then not Q(X).Null_Body then
568         Exceptional_Complete_Rendezvous(Current_Rendezvous);
569       end if;
570     end;
571   end if;
```

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```
572   if X = No_Rendezvous then
573     if Time1 > Time2 then
574       T4;
575     else
576       T3;
577     end if;
578   end if;

579 end;
```

580 8.8 Simple Accept Statements

581 The `Accept_Call` operation is provided to implement the simple **accept** statement
582 with potentially less overhead than if the `Selective_Wait` operation were used.

583 8.8.1 RTS Responsibilities

584 `Accept_Call(E,L)` shall be equivalent to

```
585   Selective_Wait(
586       (1..1=>(E=>E, Null_Body=>False)),
587       Simple_Mode,
588       P,
589       X'Access)
```

590 where `P` and `X` are compiler-generated variables to receive the parameter block
591 pointer and chosen select index values returned by `Selective_Wait`; `X` would be
592 ignored in this case.

593 8.8.2 Compiler Notes

594 `Accept_Call` should be called for an **accept** statement with a sequence of state-
595 ments. It can also be used to implement selective wait statements where the
596 compiler-generated code can determine which alternative will be selected.

597 8.9 Trivial Accepts

598 The `Accept_Trivial` operation provides a way to implement trivial accept state-
599 ments.

600 8.9.1 RTS Responsibilities

601 `Accept_Trivial(E)` shall be equivalent to

```
602   Selective_Wait(
603       (1..1=>(E=>E, Null_Body=>True)),
604       Simple_Mode,
605       P,
```

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```
606         X'Access)
```

607 where P and X are dummies that are ignored by the compiler-generated code.

608 8.9.2 Compiler Notes

609 Accept_Trivial should be called for an **accept** statement that has no internal
610 sequence of statements.

611 8.10 Completing Rendezvous

612 The Complete_Rendezvous and Exceptional_Complete_Rendezvous operations
613 are provided to complete a non-trivial rendezvous.

614 8.10.1 RTS Responsibilities

615 Before return from Complete_Rendezvous or Exceptional_Complete_Ren-
616 dezvous, the TRTS shall perform the following actions:

617 1. Lower the dispatching priority of the accepting task, if necessary.

618 2. Propagate the exception specified by the exception-ID parameter E to the caller,
619 in the case of Exceptional_Complete_Rendezvous.

620 3. Make the caller eligible for execution.

621 4. If the call was made by an asynchronous **select** statement, abort the associ-
622 ated abortable part (see Section 5).

623 The effect of these operations is undefined unless the calling task has a rendezvous
624 under way—that is, there must have been a successful preceding Selective_Wait
625 or Accept_Call operation for which no Complete_Rendezvous or Exceptional_-
626 Complete_Rendezvous operation has yet been performed.

627 The effect of calling Complete_Rendezvous or Exceptional_Complete_Ren-
628 dezvous operation when no rendezvous is in progress is undefined.

629 8.10.2 RTS Notes

630 The Complete_Rendezvous and Exceptional_Complete_Rendezvous operations
631 require the TRTS to keep track of the current caller, so that it can be awakened and
632 so that exceptions can be propagated to it, as necessary. Since **accept** and **select**
633 statements can be nested, the TRTS has to keep track of all callers being serviced by
634 a given task, and has to restore the current caller at the start of the corresponding
635 rendezvous when Complete_Rendezvous or Exceptional_Complete_Rendezvous
636 is called. This can be done by keeping a stack of current callers for each task, for
637 example.

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8.10.3 Compiler Responsibilities

The compiler shall assure that `Complete_Rendezvous` or `Exceptional_Complete_Rendezvous` operation are called only when a nontrivial rendezvous started by the `Selective_Wait` or `Accept_Call` operation is in progress.

8.10.4 Compiler Notes

`Complete_Rendezvous` should be called by the compiler-generated code for normal completion of the rendezvous, just before exit from the **accept** body. `Exceptional_Complete_Rendezvous` should be called by the compiler-generated code (or runtime system code implementing exception propagation) when an exception is propagated out of the **accept** body. The `Ex` parameter of `Exceptional_Complete_Rendezvous` should be the ID of the exception that is to be propagated to the calling task, or `Null_Exception` if there is no such exception (this has the same effect as `Complete_Rendezvous`).

These rendezvous completion operations should not be called for rendezvous initiated by the `Accept_Trivial` operation, or for cases of a selective wait where `Null_Body` is specified as `True`, since the rendezvous will be complete on return from these calls.

8.11 Count Attribute

8.11.1 RTS Responsibilities

The `Task_Count` function shall return the number of calls queued on the specified entry in the calling task. The effect of the call is undefined unless the `E` parameter is a valid entry index for the calling task. |a

8.11.2 Compiler Responsibilities

The value of `E` shall be a valid entry index for the calling task.

8.11.3 Compiler Notes

`Task_Count` should be called to obtain the value of the `Count` attribute of a task entry. |a

8.12 Callable Attribute

8.12.1 RTS Responsibilities

The `Callable` function shall return `False` if `T` identifies a task which is completed or abnormal, or if it is `Null_Task`; otherwise, it returns `True`.

The effect of the call is undefined unless the `T` parameter is valid or `Null_Task`.

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8.12.2 Compiler Responsibilities

The compiler shall assure that the argument to `Callable` is a valid task ID or `Null_Task`.

8.12.3 Compiler Notes

The `Callable` function should be called to obtain the value of the `Callable` task attribute.

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Section 9: Protected Objects

The TRTS operations required to implement protected objects are provided by the `System.Protected_Objects` package.

9.1 Description

```
with System.Compiler_Exceptions;
package System.Protected_Objects is
  Null_Protected_Entry: constant:= 0;
  Max_Protected_Entry: constant:= implementation-defined;
  type Protected_Entry_Index is range
      Null_Protected_Entry..Max_Protected_Entry;
  subtype Positive_Protected_Entry_Index is
      Protected_Entry_Index range <1..Protected_Entry_Index'Last>;
  type Call_Modes is (Simple_Call, Conditional_Call, Asynchronous_Call);
  type Protection(Num_Entries: Protected_Entry_Index) is private;
  type Protection_Access is access Protection;
  type Communication_Block is private;
  procedure Initialize_Protection(
      Object: Protection_Access;
      Ceiling_Priority: Integer);
  procedure Finalize_Protection(Object: Protection_Access);
  procedure Lock(Object: Protection_Access);
  procedure Lock_Read_Only(Object: Protection_Access);
  procedure Unlock(Object: Protection_Access);
  procedure Protected_Entry_Call(
      Object: Protection_Access;
      E: Protected_Entry_Index;
      Uninterpreted_Data: System.Address;
      Mode: Call_Modes;
      Block: out Communications_Block);
  type Barrier_Vector is
      array(Positive_Protected_Entry_Index range <>) of Boolean;
   procedure Wait_For_Completion
     (Call_Cancelled: out Boolean;
      Block: in out Communication_Block);
  procedure Wait_Until_Abortable(
      Block: in out Communication_Block);
   procedure Cancel_Protected_Entry_Call
     (Call_Cancelled: out Boolean;
      Block: in out Communication_Block);
  procedure Next_Entry_Call(
      Object : Protection_Access;
      Barriers: Barrier_Vector;
      Uninterpreted_Data: out System.Address;
      E: out Protected_Entry_Index);
  procedure Complete_Entry_Body(
```

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```
46        Object: Protection_Access;
47        Pending_Serviced: out Boolean);
48    procedure Exceptional_Complete_Entry_Body(
49        Object: Protection_Access;
50        Pending_Serviced: out Boolean;
51        Ex: Compiler_Exceptions.Exception_ID);
52    procedure Requeue_Protected_Entry(
53        Old_Object, New_Object: Protection_Access;
54        E: Protected_Entry_Index;
55        With_Abort: Boolean);
56    procedure Requeue_Task_To_Protected_Entry(
57        New_Object: Protection_Access;
58        E: Protected_Entry_Index;
59        With_Abort: Boolean);
60    function Protected_Count(
61        Object: Protection;
62        E: Protected_Entry_Index) return Natural;
63    procedure Raise_Pending_Exception(Block: Communications_Block);
64    procedure Broadcast_Program_Error;
65  private
66    -- TRTS-defined
67  end System.Protected_Objects;
```

a

68 The description of this package is divided into several subsections.

69 9.2 Representation of Protected Objects

70 Protected objects are identified to the TRTS as objects of type `Protection`.

71 9.2.1 RTS Responsibilities

72 The TRTS shall maintain data structures for controlling access to a protected object
73 and for queuing tasks on protected entries.

74 9.2.2 RTS Notes

75 The data structures used to implement protected actions are accessed only by the
76 TRTS. The data is not assumed to be organized in any particular way, and is not
77 assumed to have any particular relationship to the corresponding `Protection` ob-
78 ject. For example, the `Protection` type might contain all necessary structures, or
79 it might refer indirectly to some structures built by the TRTS. In the most extreme
80 case, the `Protection` type might consist of just its discriminant plus a pointer or
81 index to the information needed to control the corresponding protected object.

82 In general, a protected entry call may be executed by a task other than the one that
83 made the call. For this reason, an implementation may have to pass references to
84 a protected object between tasks. To allow this, most operations on `Protection`
85 objects take a parameter of `Protection_Access`. Note that this is difficult to
86 do by passing the object itself and using the `Unchecked_Access` attribute in the
87 implementation. This requires the parameter to be passed by reference, and the
88 only way to require pass by reference is to make `Protection` a tagged type. This is

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considered to be too much of a burden on the implementation.

9.2.3 Compiler Notes

For each protected object, the compiler should declare an object of the `Protection` type, constrained by the maximum entry index, at the point where the protected object declaration appears.

A protected object can contain private components declared in the corresponding protected type. These components are accessed only by compiler-generated code; the TRTS has no access to this information. The compiler is responsible for the association between this information and the corresponding `Protection` object.

9.3 Entry Indices

As with task entries, entries in protected objects are represented by nonnegative integers, with each member of an entry family considered to be an entry. However, the two entry mechanisms are different, and the range of entry values supported by the implementation of one may not be suitable for the implementation of the other. A separate type `Protected_Entry_Index` is therefore defined in this package to represent protected entries.

9.3.1 RTS Responsibilities

An entry index value `i` is *valid for a protected object* `P` if `i` is in the range `1..Num_Entries`, where `Num_Entries` is the number of entries specified by the discriminant constraint with which the corresponding `Protection` type was declared.

9.3.2 Compiler Responsibilities

The largest entry value used for a protected object shall be used as the `Num_Entries` discriminant constraint in the elaboration of the corresponding `Protection` object.

9.3.3 Compiler Notes

The compiler should assign a unique positive value of the type `Protected_Entry_Index` to each entry of a protected object. In assigning indices to entries, each component of an entry family shall be considered to be an entry. The indices should be assigned in textual order to allow for queuing policies (e.g. `Priority_Queuing`) where textual order is significant.

The assignment of entry index values to protected entries is analogous to the assignment of entry index values to task entries, as outlined in Section 8.2.3. That is, the compiler should assign a compact range of index values to protected entries, and members of entry families should be mapped onto indices as if they were arrays of entries.

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9.4 Parameter Passing

Entry bodies can be executed by tasks other than that which made the corresponding entry call. It must therefore be possible to pass entry parameters as uninterpreted data through the TRTS between tasks, as described in Section 1.9.

9.4.1 Compiler Notes

The compiler can translate the formal parameter specification of the entry to an equivalent record type, with a component for each parameter. The rendezvous parameter would then be the address of a record of this type. For parameters passed by value, the component is of the same type as the parameter and contains the parameter value. For parameters passed by reference, the component is a pointer to the actual parameter object. For parameters passed by result, the component may contain both a value and a pointer.

9.5 Initializing and Finalizing Protected Objects

9.5.1 RTS Responsibilities

Before any of the other operations described in this section can be performed on a `Protection` object, it must be initialized by calling `Initialize_Protection`. `Initialize_Protection` shall assign the specified ceiling priority to the object.

When no further calls to operations defined in this section will be made on a protected object, it should be finalized by calling `Finalize_Protection`. `Finalize_Protection` will remove any remaining calls on the entry queues of the object and raise `Program_Error` in the callers at the point of the call.

A `Protection` object is *valid* if `Initialize_Protection` has been called for it, and if `Finalize_Protection` has not yet been called for it.

9.5.2 Compiler Responsibilities

The compiler-generated code shall call `Initialize_Protection` on an object of type `Protection` before it calls any other TRTS operation with this object as a parameter. Similarly, it shall make no such calls on the object after `Finalize_Protection` is called on it. It shall call `Finalize_Protection` on a `Protection` object before deallocating it.

9.5.3 Compiler Notes

The compiler-generated code should call `Initialize_Protection` on a `Protection` object immediately after its elaboration. The value of the `Ceiling_Priority` parameter should be that specified by the `Priority` or `Interrupt_Priority` pragma for that object. If there is no such pragma, the value `System.Priority'Last` should be used, unless there is an `Interrupt_Handler` or

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`Attach_Handler` pragma appears in the corresponding protected definition, in which case `Unspecified_Priority` should be used.

The compiler-generated code should call `Finalize_Protection` on a `Protection` object immediately before any other finalization of the associated protected object.

9.6 Initiating and Completing Protected Actions

9.6.1 RTS Responsibilities

The `Lock` procedure functions as a write lock; once a task has successfully executed this call on a protected object, no other tasks will be able to complete a call to `Lock` or `Lock_Read_Only` on the same object until `Unlock` is called with the same protected object as its argument. Note that some other GNARLI operations (`Cancel_Protected_Entry_Call`, `Wait_For_Protected_Entry_Call`) may lock an object.

The `Lock_Read_Only` functions as a read lock; this procedure can be called any number of times on the same protected object. However, tasks executing `Lock` will not be able to proceed until a call to `Unlock` is made corresponding to each call to `Lock_Read_Only`.

Between a call to `Lock` or `Lock_Read_Only` and the corresponding `Unlock`, the calling task is said to be executing a *protected action* on the specified `Protection` object. If priority scheduling is in effect, a task's effective priority is raised to at least that of the protected object's ceiling priority during the protected action. A protected action should always take place with abortion deferred (see Section 5.4); the effect of calling `Lock` or `Lock_Read_Only` while abortion is not deferred is undefined.

The argument to `Unlock` must be the same as that of the most recently locked object, that is, unlocking must be in LIFO order. The result of calling `Unlock` with any other argument is undefined.

The effect of calling any of these operations with an invalid argument is undefined.

9.6.2 RTS Notes

Note that all of these procedures require that the protected object to which they apply be provided as an argument. The TRTS could provide this value for `Unlock`, since it must be the argument to the most recent call to `Lock` or `Lock_Read_Only` made by the calling task. However, since protected actions can be nested, this would require the TRTS to associate a dynamic stack of protected objects with each task. Since these operations should be as lightweight as possible, the burden of keeping track of protected action nesting is placed on the compiler-generated code, which may be able to avoid such a stack under some circumstances.

On some platforms, reader/writer locks may be much more expensive than simple mutual exclusion locking. In such cases, it may be desirable to implement `Lock` and `Lock_Read_Only` identically, allowing only one task to perform a protected action at a time, regardless of the operation being performed. This is particularly true of

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uniprocessor implementations, where the increased potential for parallelism is of no value.

9.6.3 Compiler Responsibilities

A task shall call `Unlock` only on the protected object which it most recently locked. It shall only call `Lock`, `Lock_Read_Only`, and `Unlock` on valid `Protection` objects, and with abortion deferred.

9.6.4 Compiler Notes

Before executing the code of a protected subprogram, the compiler-generated code should call `Lock` (procedure) or `Lock_Read_Only` (function) on the corresponding protected object. `Lock` should also be called before initiating a protected entry call (see Section 9.7). After the protected action is complete, `Unlock` should be called. Before completion of a protected procedure or entry call, the compiler-generated code should execute the entry bodies corresponding to entry calls on that protected object for which the barrier expression is `True`, as described in Section 9.8.3 below, only calling `Unlock` when all remaining entries are queued on `False` barriers.

RTS operations which may block should not be called while executing protected actions. This is defined as a bounded error in the Ada 9X specification. The compiler may detect and reject a compilation containing a potentially blocking operation within a protected action, but this is in general impossible. The RTS may raise `Program_Error` exception on detecting the execution of a such a potentially blocking operation; otherwise, this situation may eventually lead to deadlock.

The potentially blocking TRTS operations are:

```
Selective_Wait              Create_Task
Accept_Call                 Activate_Tasks
Accept_Trivial              Expunge_Unactivated_Tasks
Task_Entry_Call             Complete_Activation
Call_Simple                 Complete_Task
Complete_Rendezvous         Cancel_Task_Entry_Call
Abort_Tasks                 Cancel_Protected_Entry_Call
Enter_Master                Wait_For_Completion
Complete_Master             Wait_Until_Abortable
```

Other parts of the RTS (e.g. I/O) may define other potentially blocking operations that should not be called during a protected action.

This assumes that all entry calls are potentially blocking, even asynchronous select statements. This is in keeping with RM9X 4.0 9.5.1(9).

9.7 Pending Entry Calls

`Protected_Entry_Call` pends a protected entry call on the specified protected

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234 object.

235 9.7.1 RTS Responsibilities

236 `Protected_Entry_Call` shall pend an entry call on a protected object. This call will
237 either be queued or its index returned by the next call to `Next_Entry_Call` (below)
238 on this object, if the corresponding barrier is `False` or `True`, respectively.

239 The effect of the parameters is as follows:

240 1. The parameter `Object` specifies the target object.

241 2. The parameter `E` specifies the entry within the protected object to be called.

242 3. The parameter `Uninterpreted_Data` shall be passed to the entry body, as the
243 `Uninterpreted_Data` of the corresponding `Next_Entry_Call` operation (see
244 Section 1.9).

245 4. The parameter `Mode` specifies the type of call. It controls the action of `Next_-`
246 `Entry_Call`, `Complete_Entry_Body`, `Exceptional_Complete_Entry_Body`,
247 and `Wait_For_Completion`, as described below.

248 5. The parameter `Block` is a block of data whose contents is private to the GNARL
249 implementation.

a

250 9.7.2 RTS Notes

251 The semantics of this call are described in terms of "pending" the call, rather than
252 queuing it. The distinction is made because a call cannot be queued, and in particular
253 cannot be reflected in the `Count` attribute of its entry, if the barrier evaluates to `True`
254 when the call is made (RM9X 4.0, 9.5.3(8–11)).

255 The `Block` parameter can be used by the GNARL implementation to pass data
256 between this and other runtime system calls associated with an entry call. This can
257 eliminate the need to repeat complex operations within the GNARL.

258 9.7.3 Compiler Responsibilities

259 The compiler shall assure that the `Object` parameter is valid, and that the `E` pa-
260 rameter is a valid entry index for it.

261 9.7.4 Compiler Notes

262 The compiler-generated code should call `Protected_Entry_Call` with a `Mode` pa-
263 rameter of `Simple_Call` for a simple protected entry call, of `Conditional_Call`
264 for a conditional entry call, and `Asynchronous_Call` for an asynchronous **`select`**
265 statement. It should then terminate the protected action after servicing all eligible
266 pending or queued entry calls, as described under `Next_Entry_Call` below. This has
267 the effect of servicing or queuing the pended call. `Asynchronous_Call` is also used

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for implementing timed protected entry calls, as described in Section 8.5.5. Abortion should be deferred and the associated object locked before calling `Protected_Entry_Call`.

The `Uninterpreted_Data` argument should point to a record allowing access to the actual parameters of the call, as described in Section 9.4. Uninterpreted data is provided between `Protected_Entry_Call` and `Next_Entry_Call` since another task may execute the entry body.

9.8 Servicing Entry Calls

`Next_Entry_Call` is used to determine which entries have become eligible for execution as the result of a protected operation.

9.8.1 RTS Responsibilities

`Next_Entry_Call` shall return the next entry call to be serviced. The effect of calling this procedure is undefined unless the calling task is executing a protected action.

The TRTS actions for `Next_Entry_Call` are:

1. If there is a call pending on the object for which the caller is executing a protected action, and the corresponding barrier evaluates to `True`, return the entry index for that call in `E`; otherwise
2. If an element of `Barrier_Vector` corresponding to one or more awaited entries is `True`, return the entry index of one of these in `E`. Which of these is selected will depend on the entry queuing policy implemented by the TRTS. If the call is queued non-abortably and the caller has been waiting to execute an abortable part (using `Wait_Until_Abortable`, below), wake up the caller.
3. Remove the pending entry call, if any. If it was pended with `Conditional_Call` and is to be queued abortably (see Section 9.15), discard the call; otherwise, put it in the appropriate queue. This will increment the `Count` attribute for the corresponding entry.
4. If no entry call was selected, return with `E=Null_Protected_Entry`.
5. If `E` specifies an entry, `Uninterpreted_Data` will contain the value passed to `Protected_Entry_Call` in its `Uninterpreted_Data` argument when this call was made pending on the object (see Section 1.9).

Between the time that `Next_Entry_Call` returns with an entry selected and the time that the corresponding `Complete_Entry_Body` or `Exceptional_Complete_Entry_Body` is executed, a task is said to be *serving a protected entry call*.

9.8.2 Compiler Responsibilities

The compiler shall assure that `Next_Entry_Call` is called during a protected action.

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9.8.3 Compiler Notes

Immediately after a call to `Protected_Entry_Call`, and also just before unlocking the `Protection` object after leaving a protected procedure or protected entry body, the compiler-generated code should generate a vector of the current values of all barriers. This requires evaluating any barriers that might have changed. If any barrier evaluations propagates an exception, `Program_Error` should be raised in all callers of all entries of the protected object by calling `Broadcast_Program_Error`. Otherwise, `Next_Entry_Call` should be called with this vector and the `Protection` object. If a valid entry index is returned by `Next_Entry_Call` in the parameter `E`, the body of the corresponding entry should then be executed, using `Uninterpreted_Data` to access the actual parameters as described in Section 9.4.

Once locked using `Lock`, a `Protection` object should only be unlocked after a call to `Next_Entry_Call` which returns an `E` parameter of `Null_Protected_Entry`. This assumes that the only use for `Lock` is to protect some operation that changes the state of the protected object, such as the body of a protected procedure or entry [1].

9.9 Completing a Protected Entry Call

`Complete_Entry_Body` is used to inform the TRTS of the completion of an entry body.

9.9.1 RTS Responsibilities

`Complete_Entry_Body` wakes up the task waiting on the entry (if the call was pended with `Simple_Call` or `Conditional_Call`), or aborts the abortable part (if the call was pended with `Asynchronous_Call`). `Exceptional_Complete_Entry_Body` performs the same functions and also stores the indicated exception for the task that made the entry call. This exception will be raised when the target task calls `Raise_Pending_Exception`.

These procedures also return a flag indicating whether the entry body just serviced was pended by the calling task. If this flag is `True`, the entry call was serviced immediately, the call is removed, and no further processing (such as executing an abortable part) is needed.

The effect of calling either procedure is undefined unless the calling task is executing a protected action.

9.9.2 Compiler Responsibilities

The compiler shall assure that `Complete_Entry_Body` and `Exceptional_Complete_Entry_Body` is called only during while serving an entry call.

[1] The compiler-generated code may be able to determine that there are no further entries to be serviced (e.g., all barriers evaluate to `False` and neither `Protected_Entry_Call` nor `Complete_Protected_Entry_Call` was called for this object since the last `Next_Entry_Call`), in which case it can unlock the object without calling `Next_Entry_Call`.

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9.9.3 Compiler Notes

`Complete_Entry_Body` should be called at the completion of an entry body. `Exceptional_Complete_Entry_Body` should be called in cases where the entry body raises an exception, providing the propagated exception to the TRTS in the argument `Ex`. It can also be called with an argument of `Null_Exceptions` in place of `Complete_Entry_Body`.

The `Pending_Serviced` flag indicates that the body just completed was never queued; that is, that it was the pending entry. This flag should be used to avoid additional processing associated with queued calls, such as the execution of abortable parts or waiting for the completion of a simple entry call.

9.10 Propagating Exceptions

`Broadcast_Program_Error` propagates an exception to all tasks queued on the entries of a protected object.

9.10.1 RTS Responsibilities

`Broadcast_Program_Error` shall raise `Program_Error` in all callers of the object in which the caller is currently executing a protected action, and all calls to any entry of the protected object shall be dequeued.

The effect of calling this procedure is undefined unless the calling task is executing a protected action.

9.10.2 Compiler Responsibilities

The compiler shall assure that `Broadcast_Program_Error` is called during a protected action.

9.10.3 Compiler Notes

`Broadcast_Program_Error` should be called when an exception is propagated out of the evaluation of an entry barrier. This raises `Program_Error` in all callers, as required by the Ada language. The task that calls this operation may or may not be a caller of the current object, so it must be prepared for `Broadcast_Program_Error` to raise `Program_Error` in the calling task.

9.11 Example of a Protected Object

Consider the following protected object:

```
protected type PT is
  procedure P;
  function F return Boolean;
  entry E1(Param: Integer);
```

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```
375   entry E2(1..10);
376 private
377   B: Boolean:= False;
378 end PT;

379 protected body PT is

380   procedure P is
381   begin
382     S1;
383   end P;

384   function F return Boolean is
385   begin
386     S2;
387   end F;

388   entry E1(Param: Integer) when B is
389   begin
390     S3;
391   end E1;

392   entry E2(for I in 1..10) when B or F is
393   begin
394     S4;
395   end E2;

396 end PT;

397 PO: PT;
```

398 One of several equivalent ways to structure the compiler generated code and TRTS
399 calls for the above declarations is:

```
400 type PT is new Finalization.Controlled with record
401   Object: aliased Protection(Num_Entries=> 11);
402   B: Boolean:= False;
403   end record;

404 procedure Initialize(PO: in out PT) is
405 begin
406   Initialize_Protection(PO.Object'Access);                    |a
407 end Initialize;

408 procedure Finalize(PO: in out PT) is
409 begin
410   Finalize_Protection(PO.Object'Access);                      |a
411 end Finalize;

412 type E1_Params is record
413   Param: Integer;
414 end record;
415 procedure P(PO: in out PT);
416 function F(PO: PT) return Boolean;
417 procedure Service_Entries(
418     PO: in out PT;
419     Pending_Serviced: out Boolean);
```

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```
420  ...

421  procedure P_Unprotected(PO: in out PT) is
422  begin
423    S1;
424  end P_Unprotected;

425  procedure P(PO: in out PT) is
426    PS: Boolean;
427  begin
428    Defer_Abortion;                                                    |a
429    Lock(PO.Object);                                                   |a
430    P_Unprotected(PO);
431  at end
432    Service_Entries(PO, PS);
433    Unlock(PO.Object'Access);                                          |a
434    Undefer_Abortion;                                                  |a
435  end P;

436  function F_Unprotected(PO: PT) return Boolean is
437  begin
438    S2;
439  end F_Unprotected;

440  function F(PO: PT) return Boolean is
441  begin
442    Defer_Abortion;
443    Lock_Read_Only(PO.Object);
444    return F_Unprotected(PO);                                          |a
445  at end
446    Unlock(PO.Object'Access);                                          |a
447    Undefer_Abortion;                                                  |a
448  end F;

449  procedure Service_Entries(
450      PO: in out PT;
451      Pending_Serviced: out Boolean) is                                |a
452    P: System.Address;
453    subtype PO_Entry_Index is Protected_Entry_Index
454        range Null_Protected_Entry..11;
455    Barriers: Barrier_Vector(1..11);
456    E: PO_Entry_Index;
457    PS: Boolean;                                                       |a
458    Cumulative_PS: Boolean:= False;                                    |a
459  begin
460    loop
461      begin
462        Barriers(1):= PO.B;
463        Barriers(2..11):= (when others=> PO.B or F_Unprotected(PO));
464      exception
465      when others=>
466        begin
467          Broadcast_Program_Error;
468          Pending_Serviced:= True;                                     |a
469        exception
470        when Program_Error=>
471          Unlock(PO.Object'Access);                                    |
```

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```
472         Undefer_Abortion;                                                  a
473         raise;
474       end;
475     end;
476     Next_Entry_Call(PO.Object'Access, Barriers, P, E);
477     begin
478       case E is

479       when Null_Protected_Entry=>
480         exit;

481       when 1=>
482         declare
483           Params: E1_Params;
484           for Params use P;                                                a
485           Param : Integer := Params.Param;
486         begin
487            S3;
488           Complete_Entry_Body(PS);                                         a
489         end;

490       when 2..11=>
491         declare
492           type Entry_Index_Type is range 1..10;
493           I: Entry_Index_Type:= Entry_Index_Type(E - 1);
494         begin
495           S4;
496           Complete_Entry_Body(PS);                                         a
497         end;

498       end case;
499     exception                                                              |
500     when others=>                                                          |
501       Exceptional_Complete_Entry_Body(Ex=> Current_Exception, PS);         |
502     end;                                                                   |
503     Cumulative_PS:= Cumulative_PS or PS;                                   a
504   end loop;
505   Pending_Serviced:= Cumulative_PS;                                        a
506 end Service_Entries;

507 PO: PT;
```

508 This example implements the protected type as a controlled record containing the
509 state of the object. The TRTS state is represented by a `Protection` component,
510 and the private data is represented by additional fields in the record, here `B`. Con-
511 trolled type initialization and finalization are used to assure that the `Protection`
512 component is properly initialized and finalized by the TRTS.

513 Each protected subprogram takes the associated protected object as an implicit pa-
514 rameter. This example implements these directly with subprograms of the same
515 name, each of which takes an implicit parameter of type `PT`. This parameter con-
516 tains all of the state of the protected object.

517 `Barriers` represents an anonymous vector of barrier values associated with the
518 object. In the case of E1, the barrier is just a single boolean variable private to the
519 protected record, and all the barrier evaluation in `Service_Entries` accomplishes

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is to copy B into the barrier vector. This is likely to be a common case, and an important compiler optimization will be to simply allocate such a barrier variable as the corresponding element in the barrier array. In this way, no evaluation would be required; the correct value would already be in the barrier vector.

The identifier Current_Exception stands for the current exception, which the compiler generated code is responsible for knowing, possibly with the help of the exception handling portion of the RTS (The interface between the exception handler code and the exception handling portion of the RTS is not part of GNARLI).

The implementation of the E1 entry body provides an example of how entry parameters might be passed, using the Uninterpreted_Data parameter of Next_Entry_Call. The example provides a record type E1_Params with a field corresponding to each parameter of E1. The case alternative within Service_Entries implementing this body assumes that the P parameter returned by Next_Entry_Call is the address of a record of this type, with the actual parameter value in the Param field. It declares an E1_Params object at that address and places the value of this field in a local variable Params for use by the code sequence S3.

The calls to Complete_Entry_Body and Exceptional_Complete_Entry_Body return PS, a boolean flag indicating when the body just completed had never been queued, that is, that it was the pending entry. If the pending entry is completed, Service_Entries must report this using its Pending_Serviced flag, so that the caller can avoid further processing that would be required for a queued call.

The procedure P and function F are then called as normal Ada subprograms. Note that there are unprotected versions of these subprograms for use within the protected object itself; an example of the use of these unprotected subprograms can be seen in the evaluation of the barriers to E2.

Calling the entry E1 is more complex, and is discussed in detail below.

9.12 Waiting for a Protected Entry Call

Wait_For_Completion is used to wait for an entry call to complete or be cancelled.

9.12.1 RTS Responsibilities

Wait_For_Completion shall remove the protected entry call call made by the most recent call to Protected_Entry_Call executed by the calling task. If the call is queued, Wait_For_Completion blocks the calling task until the call is dequeued or the calling task is aborted.

The behavior of Wait_For_Completion varies depending on the call made:

1. Simple_Call: If the call is queued and subsequently aborted, the object is locked and the call is cancelled (dequeued). Abortion must be deferred when Wait_For_Completion is called if it is to return on abortion.

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2. `Conditional_Call`: The call will only be queued if it was requeued without abort. `Wait_For_Completion` will block until the call is complete (`Cancelled=False`) or has been requeued non-abortably (`Cancelled=True`). Since the call cannot be aborted while queued non-abortably, `Wait_For_Completion` will never dequeue the call itself, and will never lock the object.

If `Wait_For_Completion` cancels the call, `Call_Cancelled` will be `True`; otherwise, it will be `False`. Note that the object will be locked on return from `Cancel_Protected_Entry_Call` if and only if `Call_Cancelled` is `True`.

The effect is undefined if the calling task has not called `Protected_Entry_call` without an intervening call to `Wait_For_Completion`.

9.12.2 Compiler Notes

`Wait_For_Completion` is used to implement simple and conditional entry calls on protected objects. It should be called following a call to `Protected_Entry_Call` with where `Mode=Simple_Call`, after a call to to the object's `Service_Entries` procedure to execute or queue the call, and after the object has been unlocked. `Wait_For_Completion` blocks the calling task until the entry call is complete.

Entry calls may be dequeued by `Wait_For_Completion`. Dequeuing an entry call is a protected operation, and can effect the state of the protected object. The requires the object to be locked, and for any entry calls waiting on barriers that become true as a result of the dequeue to be serviced (e.g. using a `Service_Entry` routine such as that outlined in Section 9.11). To increase efficiency in cases where dequeuing is not required, the object is locked conditionally, and a flag is returned indicating whether the dequeuing was done or not. If not, the lock, entry service, and unlock operations are avoided.

The compiler should defer abortion during the call to `Wait_For_Completion` and the subsequent entry service and `Unlock` operations, if any.

9.12.3 Examples of Protected Entry Calls

A simple protected entry call such as:

```
PO.E1(1);
```

might be translated as:

```
declare
  P: E1_Params:= (Param=> 1);
  PS: Boolean;
  Cancelled: Boolean;
  B: Communication_Block;
begin
  Defer_Abortion;
  Lock(PO.Object'Access);
  Protected_Entry_Call(
      Object=> PO.Object'Access,
```

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```
597         E=> 1,
598         Uninterpreted_Data=> P'Address,
599         Mode=> Simple_Call,
600         Block=> B);
601     Service_Entries(PO, PS);
602     Unlock(PO.Object'Access);
603     if not PS then
604       Wait_For_Completion(Cancelled, B);
605       if Cancelled then
606         Service_Entries(PO, PS);
607         Unlock(PO.Object'Access);
608       end if;
609     end;
610     Undefer_Abortion;
611     Raise_Pending_Exception(B);
612   end;
```

613 Protected_Entry_Call pends the call on PO. It will be either executed or queued by
614 the call to Next_Entry_Call within Service_Entries. If the call is queued, Wait_-
615 For_Completion will wait until this call has completed by proxy or until the caller
616 is aborted. In the latter case, Wait_For_Completion locks the object, dequeues the
617 call and returns Cancelled=True, in which case entries must be serviced again.

618 Conditional entry calls are very similar, except that the Pending_Serviced flag
619 is used to determine which alternative was selected. For example, the following
620 conditional call on E2:

```
621   select
622     PO.E2(3);
623     S6;
624   else
625     S7;
626   end select;
```

627 might be translated as:

```
628   declare
629     P: E1_Params:= (Param=> 1);
630     PS: Boolean;
631     Cancelled: Boolean;
632     B: Communication_Block;
633   begin
634     Defer_Abortion;
635     Lock(PO.Object'Access);
636     Protected_Entry_Call(
637         Object=> PO.Object'Access,
638         E=> 1,
639         Uninterpreted_Data=> P'Address;
640         Mode=> Simple_Call
641         Block=> B);
642     Service_Entries(PO, PS);
643     Unlock(PO.Object'Access);
644     if not PS then
645       Wait_For_Completion(Cancelled, B);
646       PS:= not Cancelled;
647     end;
```

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```
   Undefer_Abortion;
   Raise_Pending_Exception(B);
   if PS then
     S6;
   else
     S7;
   end if;
end;
```

The mode for `Protected_Entry_Call` is `Conditional_Call`, which prevents the call to `Next_Entry_Call` in `Service_Entries` from queuing the call. If it cannot be serviced immediately, it is discarded, and `PS` will be `False` on return. `Wait_For_Completion` must still be called in case the call is requeued non-abortably. It will wait until the call is either complete or has been discarded as the result of an attempt to requeue it with abort.

9.13 Waiting for a Call to be Queued Abortably

`Wait_Until_Abortable` is used to wait for an entry call to complete or be queued abortably.

9.13.1 RTS Responsibilities

`Wait_Until_Abortable` blocks the calling task until a call is either queued abortably or until the corresponding entry body has started execution.

9.13.2 Compiler Responsibilities

`Wait_Until_Abortable` should only be called with abortion deferred.

9.13.3 Compiler Notes

`Wait_Until_Abortable` is used in the implementation of asynchronous protected entry calls. If a call has been queued abortably, the abortable part of an asynchronous **`select`** statement should not be started.

9.14 Cancelling Protected Entry Calls

`Cancel_Protected_Entry_Call` is used to removed a pending asynchronous protected entry call.

9.14.1 RTS Responsibilities

`Cancel_Protected_Entry_Call` shall remove the protected entry call made by the most recent call to `Protected_Entry_Call` executed by the calling task.

1. If the call is queued, the object is locked and the call is cancelled.

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2. If the call is in progress, `Cancel_Protected_Entry_Call` waits for it to complete.

If `Cancel_Protected_Entry_Call` cancels the call, `Call_Cancelled` will be `True` on return; otherwise, it will be `False`. Note that the object will be locked on return from `Cancel_Protected_Entry_Call` if and only if `Call_Cancelled` is `True`.

9.14.2 RTS Notes

Removing the entry call is distinguished from dequeuing the call. When an entry call is serviced, it is dequeued; any task can dequeue the call. When the call construct is complete, the call is removed; only the task that made the call can do this, and it can remove only the most recently made call. This reflects the nested structure of asynchronous entry calls; the most deeply nested structure, corresponding to the most recent pending call, should complete before any outer calls.

`Cancel_Call` always removes the most recent call, but it does not always dequeue it. It may already have been dequeued due to completion of the entry call by another task, in which case the abortable part will have been aborted.

9.14.3 Compiler Responsibilities

The compiler shall assure that `Cancel_Protected_Entry_Call` is called only when the calling task has an outstanding asynchronous protected entry call. Note that `Complete_Entry_Body` and `Exceptional_Complete_Entry_Body` may remove an asynchronous call; `Cancel_Protected_Entry_Call` should not be called in such cases.

9.14.4 Compiler Notes

The compiler should cause `Cancel_Protected_Entry_Call` to be called when the abortable part of an asynchronous **`select`** statement has completed, whether normally, by exception propagation, or by abortion. It should only be called if the calling task did not complete the entry call immediately; otherwise, it will have been removed by `Complete_Entry_Body` or `Exceptional_Complete_Entry_Body`.

As with `Wait_For_Completion`, `Cancel_Protected_Entry_Call` may dequeue an entry call, which is a protected operation. If this is done, the object will be locked on return, and the `Call_Cancelled` flag will be `True`. In this case, the object's entry calls must be serviced and the object unlocked.

The compiler should defer abortion during the call to `Cancel_Protected_Entry_Call` and the subsequent entry service and `Unlock` operations, if any.

9.14.5 Example of an Asynchronous Protected Entry Call

Consider an asynchronous call on the protected object declared in Section 9.11:

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```
select
  PO.E1(1);
  S6;
then abort
  S7;
end select;
```

This call might be translated as:

```
declare
  Cancelled: Boolean:= False;
  B: Communications_Block;
begin
  declare
    P: E1_Params:= (Param=> 1);
    PS: Boolean:= False;
  begin
    Defer_Abortion;
    Lock(PO.Object'Access);
    Protected_Entry_Call(
        Object=> PO.Object'Access,
        E=> 1,
        Uninterpreted_Data=> P'Address,
        Mode=> Asynchronous_Call,
        Block=> B);
    PO.Service_Entries(PO, PS);
    Unlock(PO.Object'Access);

    if not PS then
      Wait_Until_Abortable(B);
      Undefer_Abortion;
      Raise_Pending_Exception(B);
      S7;
    else
      Undefer_Abortion;
      Raise_Pending_Exception(B);
    end if;

  at end
    if not PS then
      Cancel_Protected_Entry_Call(Cancelled, B);
      if Cancelled then
        Service_Entries(PO, PS);
        Unlock(PO.Object'Access);
      end if;
    end if;
  end;
  Raise_Pending_Exception(B);

  if not Cancelled then
    S6;
  end if;

end;
```

The abortable part `S7` is only started if the call has been queued abortably. If the call to `Service_Entries` immediately following the call to `Protected_Entry_Call` has

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already executed the entry call, it will return with `PS=True`, and the abortable part will be avoided. If the call has been queued non-abortably, `Wait_Until_Abortable` will wait until it is queued abortably or completed. In the latter case the task will be aborted, which also prevents the abortable part from executing. The call can be queued non-abortably if the entry starts executing and is then requeued without abort.

If the entry barrier later opens and the entry body completes, the task will be marked with the nesting level of this asynchronous **select** (if not already marked for abortion to a lower nesting level) and an `Abort_Signal` pseudo-exception will be raised in the abortable part by `Complete_Entry_Body`.

The **at end** handler also checks the `PS` flag to determine whether the call was serviced immediately. If not, it calls `Cancel_Protected_Entry_Call` to remove the call. If the call is still queued, this locks the object and dequeues the call, returning with `Cancelled=True`. In this case, the object's entry barriers must be checked by `Service_Entries`.

Other details are as for asynchronous **select** with a task entry call trigger (Section 8.5.4).

9.14.6 Example of a Timed Protected Entry Call

The asynchronous form of `Protected_Entry_Call` can also be used to implement a timed entry call on a protected object. This can be done using Ada constructs directly, which are then expanded into GNARLI calls as described elsewhere in this document. For example, a timed entry call such as:

```
select
  PO.E1(1);
  S1;
or delay 10.0;
  S2;
end select;
```

might be translated as:

```
declare
  Timed_Out: Boolean:= True;
  Deadline: System.Real_Time.Time:=
      System.Real_Time.Clock + System.Real_Time.To_Time_Span(10.0);
begin
  select
    PO.E1(1);
    Timed_Out:= False;
  then abort
    delay until Deadline;
  end select;
end;
if Timed_Out then
  S2;
end if;
```

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The relative delay is translated to an absolute time, here a value of `System.Real_Time.Time`, since the select statement might be delayed an arbitrary length of time before it can attempt the protected entry call. If a relative delay were used, the actual delay would include the time required to pend the entry call, resulting in an extremely inaccurate timeout.

9.15 Requeuing an Entry Call

The `Requeue_Protected_Entry` operation is provided to requeue a protected entry call on another protected entry. `Requeue_Task_To_Protected_Entry` requeues a task entry call on a protected entry.

9.15.1 RTS Responsibilities

`Requeue_Protected_Entry` and `Requeue_Task_To_Protected_Entry` pend or queue the new entry call and return. The effects of these operations is undefined if the calling task is not serving an entry call of the appropriate type. `Requeue_Protected_Entry` takes the object defining the current entry call; `Requeue_Task_To_Protected_Entry` must determine this from the state of the GNARL implementation.

a

1. The `New_Object` parameter specifies the protected object containing the entry to be called.

2. The parameter `E` is the index of the entry on which the call is to be requeued.

3. The parameter `With_Abort` specifies that the requeued entry is to be abortable. Otherwise, abort will be deferred until the call is complete or requeued abortably. A conditional entry call will be discarded if an attempt is made to queue it abortably.

If `Requeue_Protected_Entry` is called while serving a protected entry call, its effect is dependent on whether it is requeuing the call to the same or another protected object:

1. If the two objects are the same, the current call is added to the specified queue; otherwise

2. The current call is pended on the new object.

9.15.2 Compiler Responsibilities

The compiler shall assure that the task calling `Requeue_Protected_Entry` or `Requeue_Task_To_Protected_Entry` is serving a protected or task entry call, respectively. Unless requeuing a protected entry call to another entry of the same object, `New_Object` must be locked before making this call.

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9.15.3 Compiler Notes

These operations should be called to implement a **requeue** statement where the target is a protected entry. The `With_Abort` flag should be `True` if and only if the **requeue** contains a **with abort** clause.

After calling `Requeue_Task_To_Protected_Entry`, the compiler-generated code should unlock `New_Object` after servicing its entries as described in Section 9.8.3 and leave the **accept** statement, bypassing the rest of its sequence of statements and calling `Complete_Rendezvous`. Similarly, after calling `Requeue_Protected_Entry`, the compiler-generated code should leave both protected actions in LIFO order, bypassing the rest of the entry body. The protected entry to which the call is transferred should be either parameterless or must have a parameter profile compatible with that of the entry of the calling **accept** statement or entry body.

9.15.4 Example of the Requeue of a Protected Entry Call

Consider the following example, which contains entry bodies requeuing calls between and within protected objects:

```
protected body O1 is
...
entry E1 when B1 is
begin
  S1;
  requeue O2.E2;
  S2;
end E1;
...
end O1;

protected body O2 is
...
entry E2 when B2 is
begin
  S3;
  requeue E3 with abort;
  S4;
end E2;

entry E3 when B3 is
begin
  S5;
end E2;
...
end O2;
```

One way to translate this into GNARLI calls is:

```
type T1 is new Finalization.Controlled with record
  Object: Protection(Num_Entries=> 1);
  end record;

procedure Initialize(PO: in out T1) is
```

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```
889 begin
890   Initialize_Protection(PO.Object);
891 end Initialize;

892 procedure Finalize(PO: in out T1) is
893 begin
894   Finalize_Protection(PO.Object);
895 end Finalize;

896 O1 : T1;

897 ...

898 procedure Service_Entries(
899     PO: in out T1;
900     Pending_Serviced: out Boolean) is                                    a
901   P: System.Address;
902   subtype PO_Entry_Index is Protected_Entry_Index
903       range Null_Protected_Entry..1;
904   E: PO_Entry_Index;
905   Barriers: Barrier_Vector(1..1);
906   PS: Boolean;                                                           a
907   Cumulative_PS: Boolean:= False;                                        a
908 begin
909   loop
910     begin
911       Barriers(1):= B1;

912     -- The following handler is not actually needed; there should be
913     -- no way for barrier evaluation to raise an exception in this object.
914     -- An optimizing compiler might not generate it.

915     exception
916     when others=>
917       begin
918         Broadcast_Program_Error;
919       exception
920       when Program_Error=>
921         Unlock(PO.Object'Access);                                        a
922         Undefer_Abortion;                                              a
923         raise;
924       end;
925     end;

926     Next_Entry_Call(Barriers, P, E);
927     begin
928       case E is

929       when Null_Protected_Entry=>
930         exit;

931       when 1=>
932         S1;
933         Defer_Abortion;                                                |
934         Lock(O2.Object'Access);                                        |
935         Requeue_Protected_Entry(                                       |
936             New_Object=> O2.Object'Access,                             |
937             E=> 1,                                                     |
```

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```
938             With_Abort=> False);
939         Service_Entries(O2, PS);
940         Unlock(O2.Object'Access);
941         Undefer_Abortion;                                                    a
942         goto End_of_E1;
943         S2;
944         Complete_Entry_Body(PS);                                             a
945         <<End_of_E1>>
946         null;
947       end case;
948     exception
949     when others=>                                                            |
950       Exceptional_Complete_Entry_Body(Current_Exception, PS);                |
951     end;                                                                     |
952     Cumulative_PS:= Cumulative_PS or PS;                                     |
953   end loop;                                                                  |
954   Pending_Serviced:= Cumulative_PS;                                          a
955 end Service_Entries;

956 ...

957 type T2 is new Finalization.Controlled with record
958   Object: Protection(Num_Entries=> 1);
959   end record;

960 procedure Initialize(PO: in out T2) is
961 begin
962   Initialize_Protection(PO.Object'Access);                                   a
963 end Initialize;

964 procedure Finalize(PO: in out T2) is
965 begin
966   Finalize_Protection(PO.Object'Access);                                     a
967 end Finalize;

968 O2 : T2;

969 ...

970 procedure Service_Entries(
971     PO: in out T2
972     Pending_Serviced: out Boolean) is                                        a
973   P: System.Address;
974   subtype PO_Entry_Index is Protected_Entry_Index
975       range Null_Protected_Entry..2;
976   E: PO_Entry_Index;
977   Barriers: Barrier_Vector(1..2);
978   PS: Boolean;
979   Cumulative_PS: Boolean:= False;
980 begin
981   loop
982     begin
983       Barriers(1):= B2;
984       Barriers(2):= B3;
985     exception
986     when others=>
987       begin
988         Broadcast_Program_Error;
```

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```
989         exception
990         when Program_Error=>
991           Unlock(PO.Object'Access);
992           Undefer_Abortion;
993           raise;
994         end;
995       end;
996       Next_Entry_Call(Barriers, P, E);
997       begin
998         case E is

999         when Null_Protected_Entry=>
1000          exit;

1001        when 1=>
1002          S3;
1003          Requeue_Protected_Entry(
1004              New_Object=> PO.Object,
1005              E=> 2,
1006              With_Abort=> True);
1007          goto End_of_E2;
1008          S4;
1009          Complete_Entry_Body(PS);
1010          <<End_of_E2>>
1011          null;
1012        when 2=>
1013          S5;
1014          Complete_Entry_Body(PS);

1015        end case;
1016      exception
1017      when others=>
1018        Exceptional_Complete_Entry_Body(Current_Exception, PS);
1019      end;
1020      Cumulative_PS:= Cumulative_PS or PS;
1021    end loop;
1022    Pending_Serviced:= Cumulative_PS;
1023  end Service_Entries;
```

1024 Note that Complete_Entry_Body is only called once for a call no matter how many
1025 times it is requeued. This call may return a Pending_Serviced flag of True if
1026 this "requeued" call was never actually queued, that is, if each requeue was able to
1027 process the call immediately. If Complete_Entry_Body returns with Pending_-
1028 Serviced=True, the flag will be passed up through the nested Service_Entries
1029 calls to the original caller.

1030 The first requeue is an external call, for which the compiler generates code to enter
1031 and exit the target object. Requeue_Protected_Entry returns after pending the
1032 call on the new object. The following Service_Entries will either execute or queue
1033 the call. The second requeue is an internal call, for which the compiler omits the
1034 code to lock New_Object. Requeue_Protected_Entry queues the call to the current
1035 object.

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9.16 Count Attribute

`Protected_Count` reports on the number of tasks queued on a protected entry. |a

9.16.1 RTS Responsibilities

`Protected_Count` returns the number of tasks that are waiting on the entry specified by the entry index `E` in the object on which the caller is currently executed an protected action. Note that pending calls are not included in this count. |a

The effect of calling this function is undefined unless the calling task is executing a protected action, and `E` is valid for the corresponding protected object. |a

9.16.2 Compiler Responsibilities

The compiler shall assure that `Count` is called during a protected action.

The compiler shall assure that the `E` parameter is a valid entry index on the corresponding protected object.

9.16.3 Compiler Notes

`Count` should be called to obtain the `Count` attribute for protected entries.

9.17 Rationale

9.17.1 Protected Actions—The Eggshell Model

A protected object can contain protected procedures, functions, and entries. Tasks can call these, resulting in a *protected action*, a synchronized operation on the protected object. Protected actions are often explained in terms of the *eggshell* of the protected object. A task must enter the eggshell to perform a protected operation and exit the eggshell when the operation is complete. Under priority scheduling, the effective priority of a task in the eggshell will be that of the *ceiling priority* associated with the protected object, and abortion is deferred for tasks executing in an eggshell.

Protected procedures and entries are permitted to change the state of the protected object; tasks executing these actions must be alone in the eggshell. Tasks attempting to enter the eggshell for one of these operations will be will not be able to proceed if it is already occupied. Protected functions, on the other hand, are only permitted to read the state of the object; there can be any number of tasks executing protected functions simultaneously within the eggshell. Tasks can only proceed with a protected function if there is no task inside executing a protected procedure or entry body. Protected entry calls must enter the eggshell to evaluate the entry barrier. If it is `True`, the entry body is executed; otherwise, the call is placed on the corresponding entry queue to be serviced when the barrier becomes `True`. An important aspect of the eggshell model is that all eligible calls must be served before any other task can enter the eggshell.

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9.17.2 Asynchronous Transfer of Control

The asynchronous **select** statement allows a task to execute a *triggering* statement (a delay or entry call) in parallel with an *abortable part*. If the abortable part finishes before the triggering statement starts, an attempt is made to cancel the triggering statement; cancellation can fail if the the the trigger has deferred abortion. If cancellation is successful, the asynchronous **select** completed; otherwise, it is complete on completion of the trigger. If the triggering statement finishes first, the abortable part is aborted. Unless the trigger is cancelled, the triggered statements are executed, possibly after an *Asynchronous Transfer of Control* (ATC) out of an executing abortable part.

9.17.3 Execution Models

9.17.3.1 Protected Entry Calls

Two models for the execution of protected entry calls are considered. The *self-service* model has the caller of the protected entry execute the entry body, whether it is executed immediately or after being queued. If the task is queued, it will later be awakened or interrupted (in the case of an asynchronous call) and returned to the eggshell to execute the entry body.

Arranging for a specific task to reenter the eggshell is difficult in some systems. For example, in the POSIX 1003.4a Threads specification (Pthreads), a reasonable operation for implementing the eggshell is the *mutex*, a mutual exclusion object which can be locked by a thread and is thereafter owned by that thread. It is not possible to pass ownership of a mutex to a specific task; the task owning it can only unlock it, allowing other tasks to contend for it. Some of these other tasks might not be queued entry calls with `True` barriers, but tasks attempting to enter the eggshell for some other protected operation. One of these latter tasks could get into the eggshell before all of the eligible entry calls were serviced, in violation of Ada 9X semantics.

For this reason, we also consider the *proxy* model of entry body execution. In this model, a calling task executes an entry body only if the barrier is true on entry. Otherwise, the call is queued and, if not canceled, executed by the task (the proxy) that causes the barrier to become `True` before it exits the eggshell.

This interface is designed to support the proxy model of execution. Since it does not require the passing of locks between threads, it is less demanding on underlying system support for mutual exclusion, and so should be efficiently implementable on a larger number of platforms. Its major disadvantage is an unbounded execution time on exit from the eggshell of a protected object. However, this assumes an unbounded cascade of entry calls becoming eligible to execute as the result of a single protected operation, which is likely to be rare in practice.

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9.17.3.2 Asynchronous Transfer of Control

Similarly, two models for the execution of an asynchronous **select** statement are considered. The *two-thread* model spawns a separate thread of control, with one thread executing the abortable part, and the other executing the entry call. On completion of the entry call, the thread that is executing the abortable part can be cancelled. When cancelled, the abortable part's finalization actions must be performed, either by the thread executing the abortable part (via asynchronous transfer of control) or by the thread executing the triggering statement.

If the proxy model is used, the triggering statement will be executed by one of the threads only in the case where the barrier is `True` at the time of the call. Otherwise, the thread spawned to execute the triggering statement will be blocked until the entry body is executed by some other task.

The *one-thread* model either executes the triggering statment immediately, or queues the asynchronous call and then continues with the execution of the abortable part. If protected entries are implemented using the self-service model, continuation of the triggering statement (i.e. by expiration of the delay or by the barrier expression becoming `True`) interrupts the abortable part. If the proxy model is used, the triggering statement may be executed in parallel with the abortable part by another task, with the abortable part being aborted only on the completion of the triggering statement. At the completion of the triggering statement, the caller executes the triggered statements; aborting the abortable part if necessary.

The two-thread model adds weight to the asynchronous **select**, since it requires a thread with its own stack to be created and destroyed. This can be expensive in both time and space. Also, the increased parallelism of the two thread model is limited by the proxy model of protected entry execution, since one of the threads will often be blocked until the entry body has been executed by another task. Finally, the need to treat multiple threads as one task complicates many other TRTS functions such as rendezvous. For these reasons, and because supporting both models while avoiding TRTS callbacks to user code is difficult, this interface is designed to support the one-thread model.

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Section 10: Runtime Support for Storage Management

This section describes support for storage allocation provided by the TRTS.

10.1 Package Specification

```
with System.Storage_Elements;
package System.Storage_Allocation is

   procedure Allocate_Block(
       Storage_Address: out System.Address;
       Size_In_Storage_Elements: System.Storage_Elements.Storage_Count;
       Alignment: in System.Storage_Elements.Storage_Count);

   procedure Deallocate_Block(Storage_Address: System.Address);

   function Maximum_Storage return System.Storage_Elements.Storage_Count;

end System.Storage_Allocation;
```

10.2 RTS Responsibilities

`Allocate_Block` shall allocate `Size_In_Storage_Elements` storage units of contiguous storage, aligned according to `Alignment`, and return the address of the first storage unit in `Storage_Address`. None of the storage in this block will be allocated by further calls to `Allocate_Block` until `Deallocate_Block` is called with the value returned in `Storage_Address`. If there are insufficient resources to allocate the requested storage, `Storage_Error` shall be raised at the point of the call.

`Deallocate_Block` shall make the storage starting at `Storage_Address` available for further allocations using `Allocate_Block`.

`Maximum_Storage` shall return the maximum storage that can be allocated at any one time. If the RTS is unable to determine this size, `Maximum_Storage` shall return zero.

`Allocate_Block`, `Deallocate_Block`, and `Maximum_Storage` shall be synchronized with respect to tasking, that is, they will function correctly when called from multiple tasks, regardless of the relative timing of the various calls.

10.3 RTS Notes

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The TRTS should use **new** to allocate all Ada objects not supplied as parameters from the compiler or allocated through elaboration of object declarations, in order to allow the compiler to perform the necessary initializations on the allocated objects. However, it should not depend on the default storage pool implementation provided by the compiler, as this storage pool may use TRTS calls, possibly leading to an infinite recursion in the TRTS. Instead, the TRTS should implement its own storage pool using the underlying operating system or machine services, and use this storage pool for all TRTS allocations. For example, this TRTS storage pool could be implemented with the operations defined in the `System.Storage_Management` package (though they need not):

```
with System.Storage_Pools;
with System.Storage_Elements;
package System.RTS_Storage_Pools is

   type RTS_Storage_Pool is new
       System.Storage_Pools.Root_Storage_Pool;

   procedure Finalize(Object: in out RTS_Storage_Pool);

   procedure Allocate(
       Pool: in out RTS_Storage_Pool;
       Storage_Address: out System.Address;
       Size_In_Storage_Elements: Storage_Elements.Storage_Count;
       Alignment: Storage_Elements.Storage_Count);

   procedure Deallocate(
       Pool: in out RTS_Storage_Pool;
       Storage_Address: System.Address;
       Size_In_Storage_Elements: Storage_Elements.Storage_Count;
       Alignment: Storage_Elements.Storage_Count);

   function Storage_Size(Pool: RTS_Storage_Pool)
       return Storage_Elements.Storage_Count;

end System.RTS_Storage_Pools;

with System.Storage_Allocation; use System.Storage_Allocation;
package body System.RTS_Storage_Pools is

   procedure Finalize(Object: in out RTS_Storage_Pool);
   begin
     null;
   end;

   procedure Allocate(
       Pool: in out RTS_Storage_Pool;
       Storage_Address: out System.Address;
       Size_In_Storage_Elements: System.Storage_Elements.Storage_Count;
       Alignment: in System.Storage_Elements.Storage_Count) is
   begin
     Allocate_Block(Storage_Address, Size_In_Storage_Elements, Alignment);
   end Allocate;

   procedure Deallocate(
```

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```
73          Pool: in out RTS_Storage_Pool;
74          Storage_Address: System.Address;
75          Size_In_Storage_Elements: System.Storage_Elements.Storage_Count;
76          Alignment: in System.Storage_Elements.Storage_Count) is
77       begin
78         Deallocate_Block(Storage_Address);
79       end Deallocate;

80       function Storage_Size(Pool: RTS_Storage_Pool)
81           return System.Storage_Elements.Storage_Count is
82       begin
83         return Maximum_Storage;
84       end Storage_Size;

85    end System.RTS_Storage_Pools;
```

86 10.4 Compiler Notes

87 The compiler is responsible for implementing the default storage pool for allocations
88 using **new**. How it does this is not specified by the TRTS. The TRTS provides a
89 synchronized allocation service for this purpose, but the compiler need not use it.
90 It may decide instead to allocate all storage from the stack, using TRTS services to
91 provide the necessary synchronization. It may decided to use `Allocate_Block` to
92 allocate all of the storage for a pool, and manage the storage within this block itself,
93 again using TRTS synchronization. It may use a combination of these methods,
94 allocating storage for some pools on the stack and others using `Allocate_Block`.

95 10.5 Rationale

96 In cases where memory allocation is not synchronized, or where it is implemented
97 over a bare machine, the TRTS can still provide synchronization at least as easily
98 as the compiler can do so by using TRTS synchronization calls. On the other hand,
99 the TRTS will sometimes be built over multithreaded operating system services
100 which provide thread-safe (i.e. synchronized) memory allocation primitives, and
101 these may be the only allocation primitives available on the platform. In this case,
102 the only service available will be synchronized, whether or not GNARLI specifies it
103 as synchronized. In these cases, it would be wasteful for the compiler to assume
104 that memory allocation needed additional synchronization. For these reasons, the
105 GNARLI interface provides synchronized memory allocation to the compiler.

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106 *This side is intentionally blank.*

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Section 11: Compiler Support for Exceptions

The `System.Compiler_Exceptions` package provides an interface to the compiler's implementation of exceptions for the TRTS implementation. The procedures in this package are not part of the TRTS. The TRTS interface makes the implementation of exceptions and exception handling entirely the responsibility of the compiler. The package is either provided by the compiler developer or written by the TRTS builder to fit the exception-handling mechanism used by the compiler.

11.1 Package Specification

```
with System.Machine_Specifics;
package System.Compiler_Exceptions is
  type Exception_ID is private;
  Null_Exception: constant Exception_ID;
  Constraint_Error: constant Exception_ID;
  Program_Error: constant Exception_ID;
  Storage_Error: constant Exception_ID;
  Tasking_Error: constant Exception_ID;
  procedure Raise_Exception(E: Exception_ID);
  procedure Notify_Exception(
      Which: Machine_Specifics.Machine_Exceptions;
      Info:  Machine_Specifics.Error_Information;
      Modified_Registers: Machine_Specifics.Pre_Call_State);
  subtype Exception_ID_String is String(1..16);
  function Image(E: Exception_ID) return Exception_ID_String;
private ... -- compiler-defined
end System.Compiler_Exceptions;
```

11.2 Compiler Responsibilities

The `Exception_ID` type shall be used to identify exceptions at the TRTS interface. The implementation of this type shall be determined by the compiler developer. For each of the Ada language-defined exceptions that may need to be raised by the TRTS, a constant of the `Exception_ID` type shall be provided in this package, that identifies that exception.

The special value `Null_Exception` shall identify no exception. The `Raise_Exception` operation shall have the effect of raising the specified exception in the calling task. Calling `Raise_Exception` with an argument of `Null_Exception` shall have no effect.

The `Notify_Exception` call shall translate a machine exception into an Ada exception, if any, raising that exception in the affected task. If the machine exception

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does not require the raising of an Ada exception, `Notify_Exception` shall have no effect. It must be called when a machine-specific condition arises that requires an Ada exception to be called. If called from within an error handler, it will raise the Ada exception at the point that the affected task was interrupted. If called from within a trap handler, this call may exit the handler, returning control to application code. This procedure has three parameters:

1. `Which` shall specify which hardware exception occurred. Its type, `Machine_Exceptions` need bear no particular relationship to `Exception_ID` or to any actual Ada exceptions. The values of this type need correspond to (only) the exceptions that are detected by hardware and need to be mapped to Ada. These values may include application-specific conditions that are detected by hardware, as well as conventional machine-specific error conditions. Some predefined Ada exceptions may not be represented (for example if the hardware does not detect overflow), or there may be several hardware exceptions (e.g, fixed-point overflow, floating-point underflow, etc.) that are mapped to a single exception (e.g., `Constraint_Error`) by the compiler's implementation of exceptions.

2. `Info`. This shall specify other machine-level information, that may be specific to the particular exception.

3. `Modified_Registers`. This shall specify the values of any registers that the TRTS needed to change from the values which they had at the time of the hardware exception in order to call `Notify_Exception`. This shall include the old program counter value at the time execution was interrupted.

The `Image` function returns a string that uniquely identifies the exception specified by the exception ID parameter, within the calling program. The exception implementation shall document a mechanism for using this string to determine the corresponding Ada exception declaration.

11.3 Compiler Notes

There are a number of ways exception handling may be implemented. They are all fairly complex, and it is not clear that there is one best way. The interface between the compiler's implementation of exceptions and the TRTS is therefore designed to hide all details of the implementation of exceptions from the TRTS. For interface purposes, all exceptions are assigned unique identifiers (e.g., integers) and the predefined exceptions are assigned standard identifiers known to the TRTS.

Ada 9X defines the package `System.Exceptions` which defines the `Exception_Occurrence` type. An Ada 9X exception handler can obtain the value of an exception occurrence, which specifies the exception that caused the handler to be invoked and other implementation-defined information related to the exception and the environment in which it was raised. This package defines the procedure `Reraise_Occurrence`, which takes an exception occurrence and raises it again.

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This interface can be used in place of `System.Compiler_Exceptions`. However, there may be problems with implementation-specific data attached to the occurrences. For example, when an exception is raised in an **accept** statement, it must be raised as well in the caller. If this exception occurrence is reraised in the caller, the information associated with it will be consistent with the environment where it was raised, which is in another task. This could be particularly disconcerting if the information contained, for example, a stack traceback.

However, `System.Exceptions` can be used as the basis for the implementation of `System.Compiler_Exceptions`. The compiler only needs to provide an implementation of `Raise_Exception` that can take an exception occurrence and raise a new occurrence of the same type.

11.4 RTS Responsibilities

The argument to `Raise_Exception` shall be a valid value of the type `Exception_ID`. This means that it will either be one of the constant values defined in `System.Compiler_Exceptions`, or a value provided by the compiler as a parameter to some GNARLI operation.

11.5 RTS Notes

This interface has been designed so as to avoid unnecessarily restricting the implementation of exceptions and exception handling, which is considered to be entirely the responsibility of the compiler. However, under certain circumstances the TRTS is required to cause an exception to be raised in an Ada task. This requires an interface between the implementation of exception handling and the TRTS. During recovery from an exception, there may be need to propagate an exception to another task, or to abandon execution of a task.

`Raise_Exception` is provided so that the TRTS can raise an exception in a task that may have been interrupted, blocked, or have called the TRTS—by calling this procedure within the context of that task. For example, it is needed for the case where the exception specified by parameter `E` of `Exceptional_Complete_Rendezvous` must be propagated to a task that is blocked as caller in a rendezvous. Note that since the value of `E` is not known until run time, the TRTS cannot use a **raise** statement to do this.

The TRTS implementation should call `Notify_Exception` when a machine exception occurs that may require raising an Ada exception. It should be called when the error occurs, without allowing the task to execute further. `Notify_Exception` should be called with arguments provided by the underlying implementation, typically arguments to an error handler.

Note that TRTS interface requires that the implementation of recovery from exceptions take care when propagating exceptions to recover storage collections that are no longer needed, to call `Complete_Master` when exiting a master, and to call `Complete_Task` when abandoning execution of a task.

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11.5.0.3 GNARL Implementation Notes

The CARTS definition of `Raise_Exception` differs from that given above in that it had a `Pre_Call_State` parameter. This was intended for use where the TRTS had responsibility for context switching. A task would be marked as needing to raise an exception the next time that it ran. When the context was switched to this task, the context switch code would restore as much of the task state as possible and put the rest into a `Pre_Call_State` object. `Raise_Exception` would then be called to restore the remainder of the state and cause the exception to be raised. An example of state that might not be restored at this point is the program counter, since changing the program counter would cause the context switch code to loose control before it would be able to call `Raise_Exception`.

In general, the GNARLI will be implemented using the GNULLI specification for low-level tasking support. This package hides the context switch from the GNARLI implementation, so there is no mechanism for interposing a call to `Raise_Exception` in the context switch. Instead, a task can use polling to determine if it should raise an exception. For this purpose, it is only necessary to raise an exception in the current context, so the GNARLI specification of `Raise_Exception` has only a single parameter specifying the exception to be raised.

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Section 12: Acknowledgements

The preparation of this report was funded in part by the Air Force, contract F08630-92-C-0019, through the GNAT project; and by the Ada Joint Program Office under the Ada Technology Insertion Program, through the U.S. Army CECOM, contract DAAB07-91-D-Q502/D.O. 0006/0017.

Portions of this document are reused from drafts of earlier runtime system interface definitions written by T.P. Baker. These are the Corset interface [3, 4], the ARTEWG Model Runtime System Interface [5], and the CARTS Software Interface Specification [2].

The MRTSI was originally drafted by T.P. Baker, then developed further by the Ada Run Time Environment Working Group (ARTEWG) of the Special Interest Group for Ada (SIGAda) of the Association of Computing Machinery (ACM), and subsequently published in *Ada Letters*. The CARTS document was written by by T.P. Baker for Westinghouse Electric Corporation and the Avionics Laboratory, Wright Laboratory, Aeronautical Systems Division, United States Air Force. Portions of the MRTSI document were reused, with permission of the ARTEWG.

The technical contributions of the ARTEWG members who worked on the MRTSI document are gratefully acknowledged, as are those of the CARTS team members who reviewed the CARTS document. The Avionics Laboratory and Westinghouse Electric Corporation are gratefully acknowledged for granting permission to reuse portions of the CARTS document.

Thanks are also due to Offer Pazy, for his many helpful comments and suggestions.

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